

The Generalized Temporal Role Based Access Control Model

Security Management
Lecture 7

March 21, 2006



Outline

- Introduction and Motivation
- Overview of the Generalized Temporal RBAC Model
- Expressiveness and Design Considerations
- Related Work
- Conclusion and Future Work

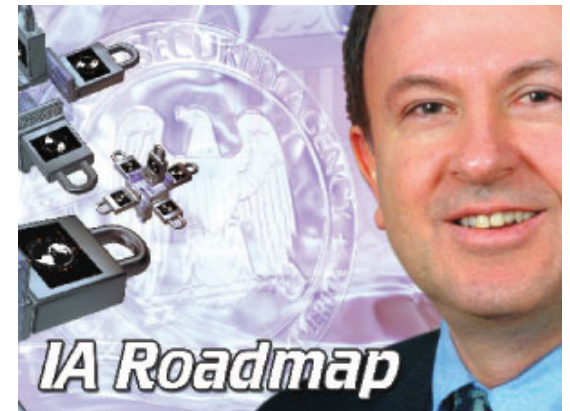


Research Motivation

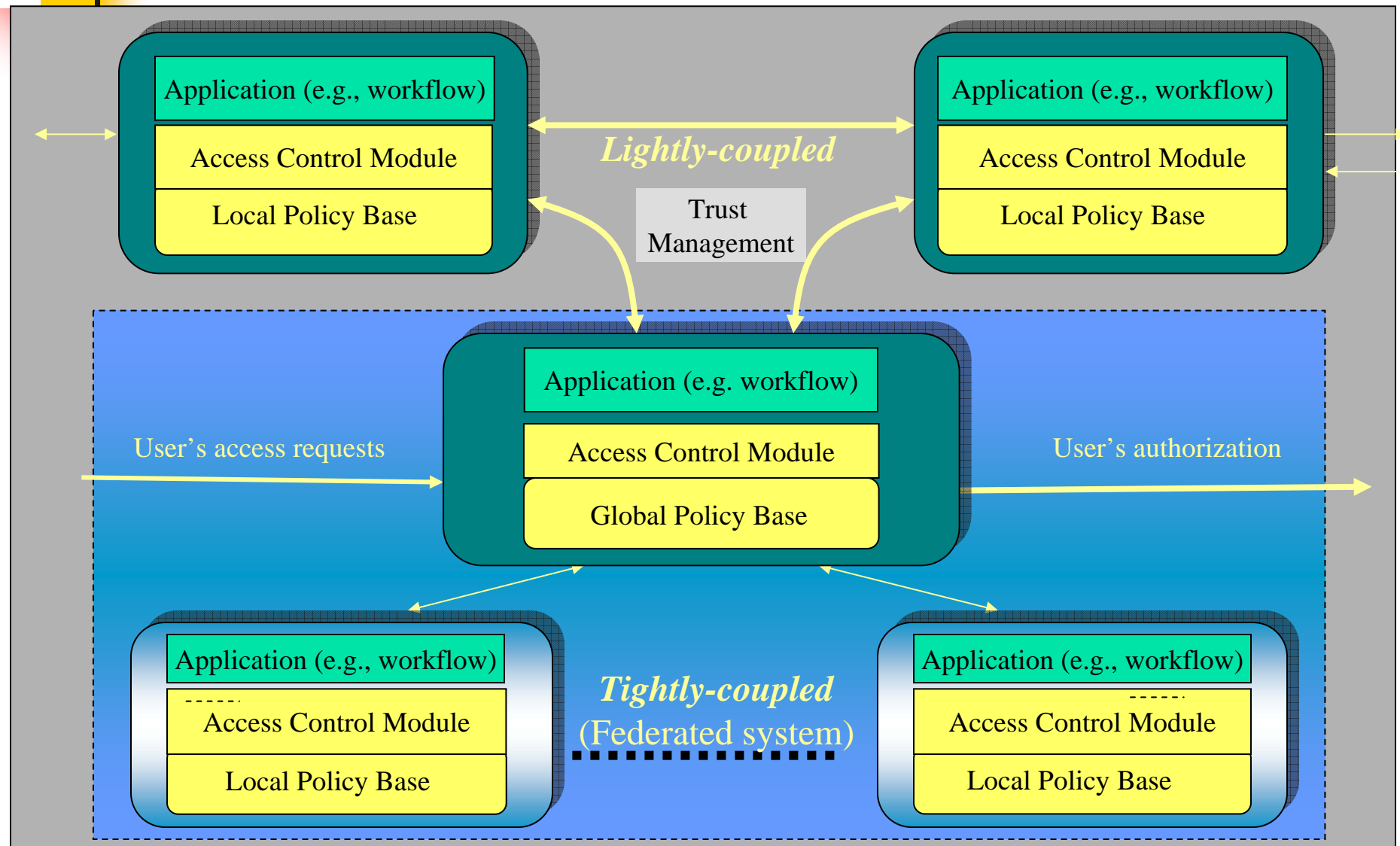
- Insider attack is a major threat in organizational systems (CSI/FBI Survey)
- Traditional discretionary and mandatory access control (DAC & MAC) approaches have limitations
- Context-based access control is a critical need for emerging applications

Research Motivation

- “To realize the Department of Defense's vision for the Global Information Grid (GIG), information assurance (IA) requirements include robust identity, authentication and **privilege management, policy for dynamic access control, security management,** and 'persistence monitoring' or continual monitoring throughout the network, according to Daniel G. Wolf, the director of information assurance for the National Security Agency (NSA).”



Security Management in Multidomain Environment

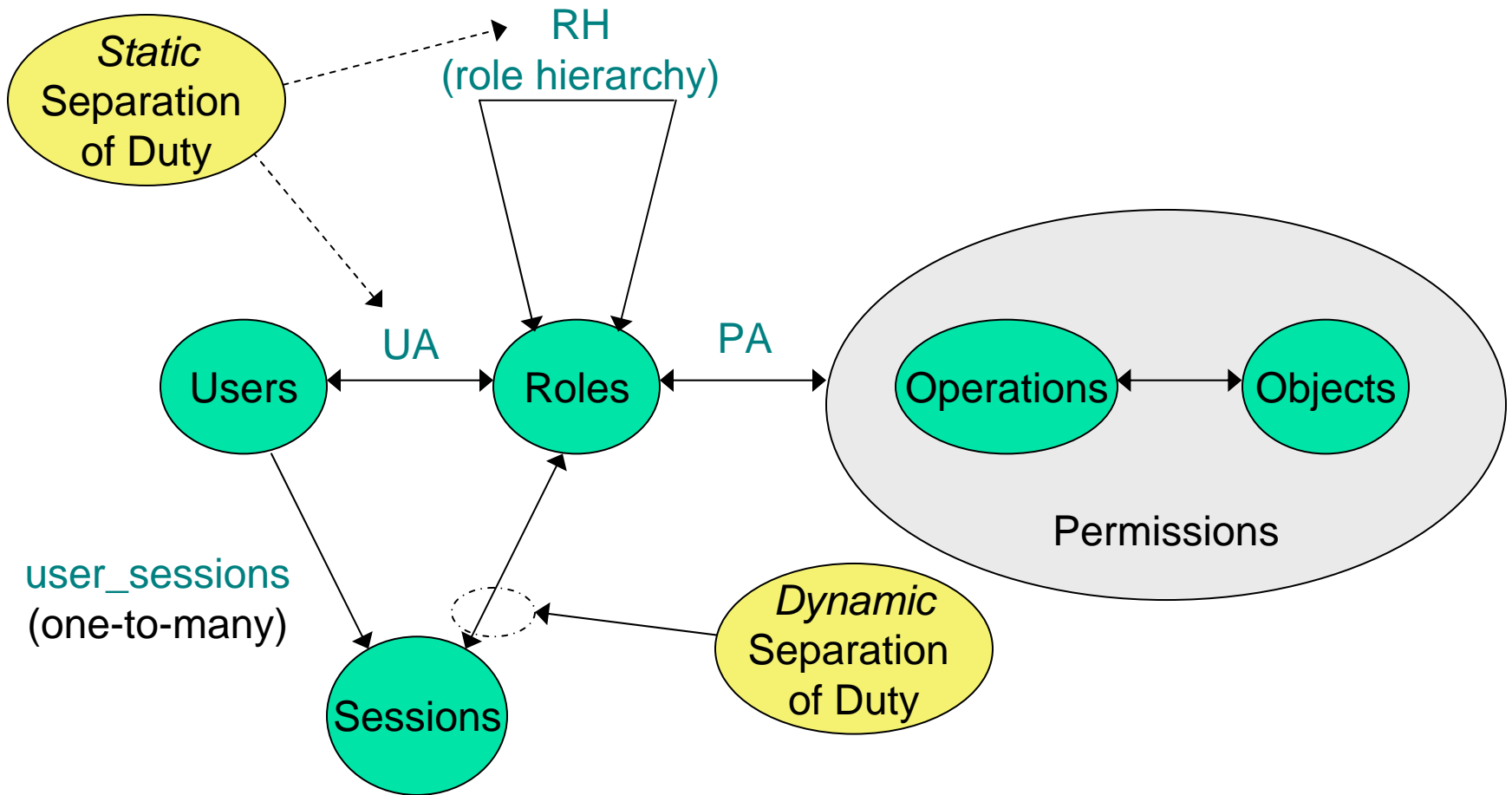




Role Based Access Control (RBAC)

- RBAC is a promising approach for addressing diverse security needs
- Access control in organizations is based on “roles that individual users take on as part of the organization”
- A role is “is a collection of permissions”
- Constraints are applied to all the links

NIST Constrained RBAC





Advantages of RBAC

- Allows efficient security management
- Supports principle of least privilege
- Separation of duty constraints
- Policy-neutral and provides generality
- Encompasses traditional discretionary and mandatory policies
- Suitable for use in multidomain environments such as digital government, multi-enterprise venture and international coalition



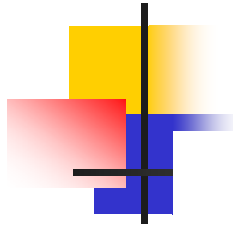
Time-based Access Control Requirement

- Organizational functions and services with temporal requirements
 - A **part-time staff** is authorized to work only between 9am-2pm on weekdays
 - A **day doctor** must be able to perform his/her duties between 8am-8pm
 - An **external auditor** needs access to organizational financial data for a period of three months
 - A video library allows access to a **subscriber** to view at most three movies every week
 - In an insurance company, an **agent** needs access to patient history until a claim has been settled

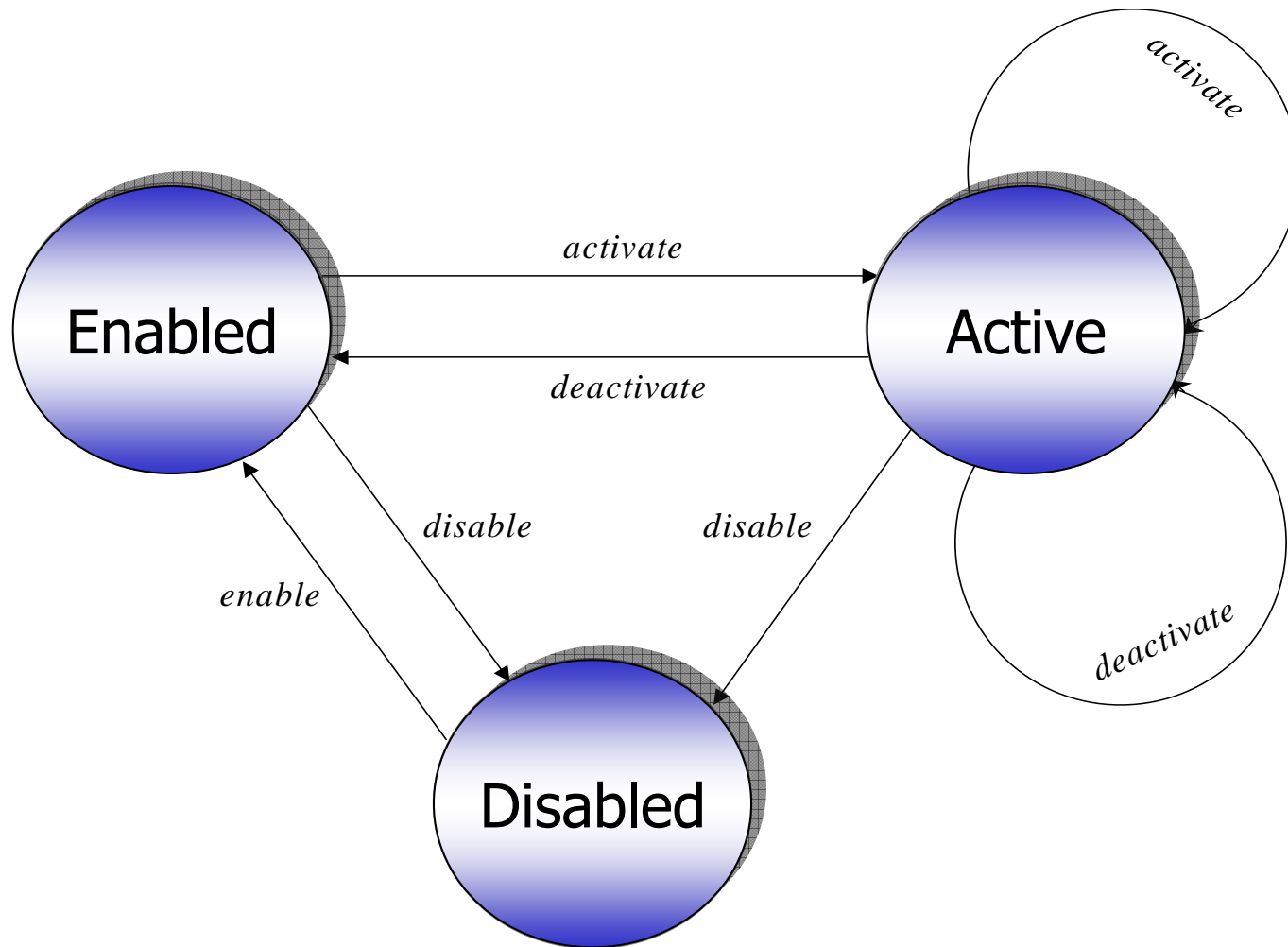


What to model in Generalized Temporal RBAC (GTRBAC)?

- Triggers and Events
- Temporal constraints
 - Roles, user-role and role-permission assignment constraints
 - Activation constraints (cardinality, active duration,..)
- Temporal role hierarchy
- Time-based Separation of duty constraints



States of a Role in GTRBAC





Temporal Constraints: Roles, User-role and Role-permission Assignments

- Periodic time
 - $(I, P) : \langle [begin, end], P \rangle$ is a set of intervals
 - P is an infinite set of recurring intervals
- Calendars:
 - *Hours, Days, Weeks, Months, Years*
- Examples
 - all.Weeks + {2, ..., 6}.Days + 10.Hours ▷ 12.hours*
 - Daytime (9am to 9pm) of working days



Temporal Constraints: Roles, User-role and Role-permission Assignments

- Periodicity: $(I, P, pr:E)$
 - $([1/1/2000, \infty], \text{Daytime}, \text{enable DayDoctor})$
 - $([1/1/2001, \infty], \{\text{Mon,Wed}\}, \text{assign}_U \text{DayDoctor to Smith})$
- Duration constraint: $(D, pr:E)$
 - $(\text{Five hours}, \text{enable DoctorInTraining})$
 - $\text{activate DayDoctor for Smith} \rightarrow \text{enable DoctorInTraining after 1 hour}$



Activation Time Constraints

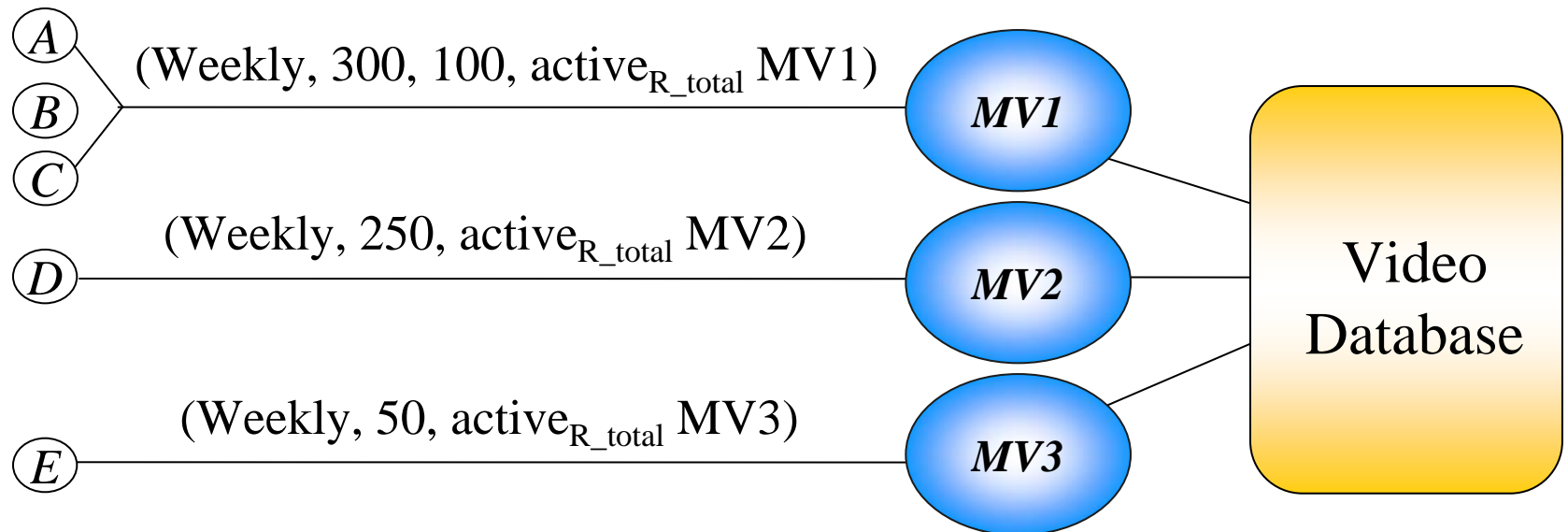
- Active role duration
 - Total duration for role activation
 1. Per role: $D_{\text{active}}, [D_{\text{default}}], \text{active}_{R_{\text{total}}} r$
 2. Per user role: $D_{\text{uactive}}, u, \text{active}_{UR_{\text{total}}} r$
 - Max active role duration per activation
 1. Per role: $D_{\text{max}}, \text{active}_{R_{\text{max}}} r$
 2. Per user role: $D_{\text{umax}}, u, \text{active}_{UR_{\text{max}}} r$
- Cardinality
 - Total number of role activations
 1. Per role: $N_{\text{active}}, [N_{\text{default}}], \text{active}_{R_{\text{n}}} r$
 2. Per user role: $N_{\text{uactive}}, u, \text{active}_{UR_{\text{n}}} r$
 - Max number of concurrent activations
 1. Per role: $N_{\text{max}}, [N_{\text{default}}], \text{active}_{R_{\text{con}}} r$
 2. Per user role: $N_{\text{umax}}, u, \text{active}_{UR_{\text{con}}} r$

Example GTRBAC access policy for a healthcare system

1	a.	(<i>DayTime</i> , enable DayDoctor), (<i>NightTime</i> , enable NightDoctor)
	b.	((M, W, F), assign _U <i>Adams</i> to DayDoctor), ((T, Th, S, Su), assign _U <i>Bill</i> to DayDoctor); ((M, W, F), assign _U <i>Alice</i> to NightDoctor), ((T, Th, S, Su), assign _U <i>Ben</i> to NightDoctor)
	c.	([10am, 3pm], assign _U <i>Carol</i> to DayDoctor)
2	a.	(assign _U <i>Ami</i> to NurseInTraining); (assign _U <i>Elizabeth</i> to DayNurse)
	b.	<i>c1</i> = (6 hours, 2 hours, enable NurseInTraining)
3	a.	(enable DayNurse → enable <i>c1</i>)
	b.	(activate DayNurse for <i>Elizabeth</i> → enable NurseInTraining after 10 min)

Example of Activation Time Constraint

- Video library offers 600 hours of total time per week
- *A*, *B* and *C* subscribe for 100 hours each
- *D* subscribes for 250 hours
- *E* subscribes for 50 hours

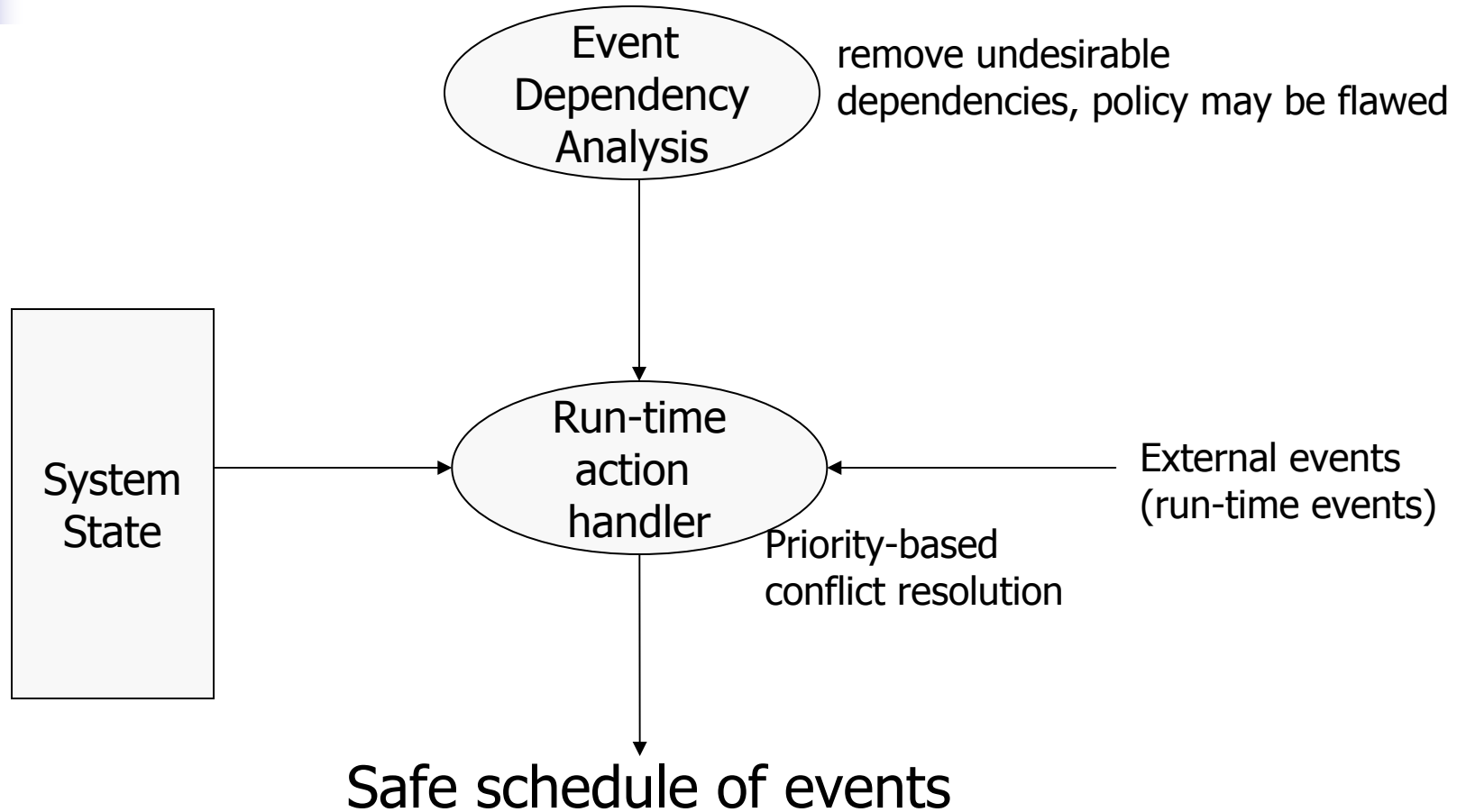




GTRBAC Execution Model for Handling Anomalies

- GTRBAC specification can encounter two possible anomalies
 - Simultaneous occurrence of conflicting events
 - Arbitrary triggering of interdependent triggers can create ambiguity
- Conflict resolution
 - Higher priority takes precedence
 - Disabling event takes precedence when the priorities of the conflicting events are the same
 - *disable r* takes precedence over *enable r*
 - More specific constraint overrides

GTRBAC Execution Model





Conflicts in GTRBAC

- GTRBAC specification can generate 3 types of conflicts
 - *Type 1*: between events of same type but opposite nature,
 - e.g., `enable r` *VS.* `disable r`
 - *Type 2*: between events of dissimilar types
 - e.g., `activate r for u` *VS.* `de-assign r to u` OR `disable r`
 - *Type 3*: between constraints
 - (a) $(X, pr:E)$ *VS.* $(X, q:E)$ OR $(X', q:Conf(E))$
 - (b) Per-role *VS.* per-user-role constraints



Handling Conflicts

- *Type 1 and Type 3(a)*
 - Higher priority takes precedence
 - Disabling event takes precedence if priorities are the same
 - e.g., disable r takes precedence over enable r
- *Type 2*
 - activation event has lower precedence
- *Type 3(b)*
 - per-user-role constraints takes precedence
- Example
 - $\{H:\text{enable } r_0, H:\text{disable } r_0, VH:\text{enable } r_1, H:\text{disable } r_1, VH:(s:\text{activate } r_1 \text{ for } U)\}$
 - After resolution
 $\{H:\text{disable } r_0, VH:\text{enable } r_1, VH:(s:\text{activate } r_1 \text{ for } U)\}$



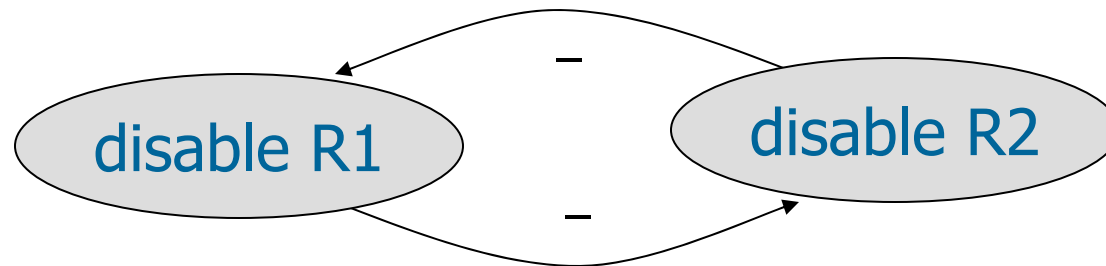
Ambiguous Event Dependency

- A set of triggers may give rise to ambiguous semantics
- Example:
 - $tr1$: enable R1 \rightarrow disable R2
 - $tr2$: enable R2 \rightarrow disable R1
 - Let the runtime requests be: {enable R1; enable R2},
 1. $tr1$ fires: {enable R1; disable R2}
(Intuitively, $tr1$ blocks $tr2$)
 2. $tr2$ fires: {enable R2; disable R1}
(Intuitively, $tr2$ blocks $tr1$)
- Solution: Detect ambiguity using Labeled dependency graph

two symmetric possibilities

Dependency Graph Analysis

- Labeled Dependency Graph
 - Directed graph (N, E)
 - N : set of prioritized events that occur in the head of some trigger
 - E : set of triples of the form (X, I, Y)
 - For all triggers $[B \rightarrow p:E]$
 - For all events E' in the body B , and for all nodes $q:E'$ in N
 - $\langle q:E', +, p:E \rangle$
 - $\langle r:\text{conf}(E'), -, p:E \rangle$ for all $[r:\text{conf}(E')]$ in N such that $q \leq r$
- Dependency Graph for the Example:





Safe Set of Triggers

- A set of triggers T is *safe* if its labeled dependency graph has no cycles with label "-".
- **Theorem:** If a T is *safe*, then there exists exactly one execution model.
- **Complexity of *DAG-based safeness algorithm*** : $O(|T|^2)$.



Role Hierarchy in GTRBAC

- Useful for efficient security management of an organization
- No previous work has addressed the effect of temporal constraints on role hierarchies
- GTRBAC-based temporal role hierarchies allow
 - Separation of permission inheritance and role activation semantics that facilitate management of access control
 - Capturing the effects of the presence of temporal constraints on hierarchically related roles



Axioms

- **Axioms:** *For all $r \in \text{Roles}$, $u \in \text{Users}$, $p \in \text{Permissions}$, $s \in \text{Sessions}$, and time instant $t \geq 0$, the following implications hold:*
 1. $p_assigned(p, r, t) \rightarrow can_be_acquired(p, r, t)$
 2. $u_assigned(u, r, t) \rightarrow can_activate(u, r, t)$
 3. $can_activate(u, r, t) \wedge can_be_acquired(p, r, t) \rightarrow can_acquire(u, p, t)$



Unrestricted Hierarchies

formal definitions

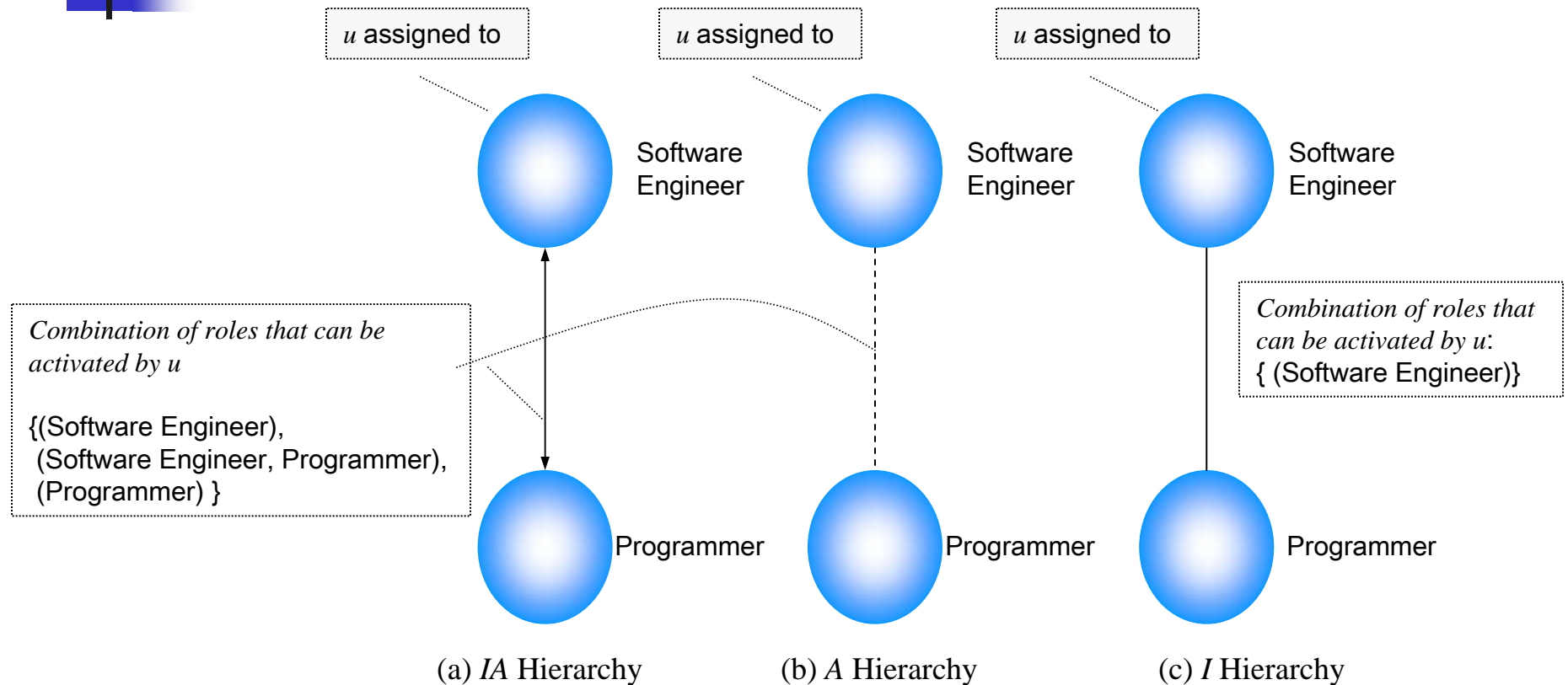
- Unrestricted *I*-hierarchy ($x \geq_t y$)
 $\forall p, (x \geq_t y) \wedge can_be_acquired(p, y, t) \rightarrow can_be_acquired(p, x, t)$
- Unrestricted *A*-hierarchy ($x \succcurlyeq_t y$)
 $\forall u, (x \succcurlyeq_t y) \wedge can_activate(u, x, t) \rightarrow can_activate(u, y, t)$
- Unrestricted *IA*-hierarchy ($x \succeq_t y$)
 $(x \succeq_t y) \rightarrow (x \geq_t y) \wedge (x \succcurlyeq_t y)$
- Consistency Property:
Let $\langle f \rangle \in \{\geq_t, \succcurlyeq_t, \succeq_t\}$ and $\langle f' \rangle \in \{\geq_t, \succcurlyeq_t, \succeq_t\} / \{\langle f \rangle\}$. Let x and y be distinct roles such that $x \langle f \rangle y$; then the condition $\neg(y \langle f' \rangle x)$ must hold.



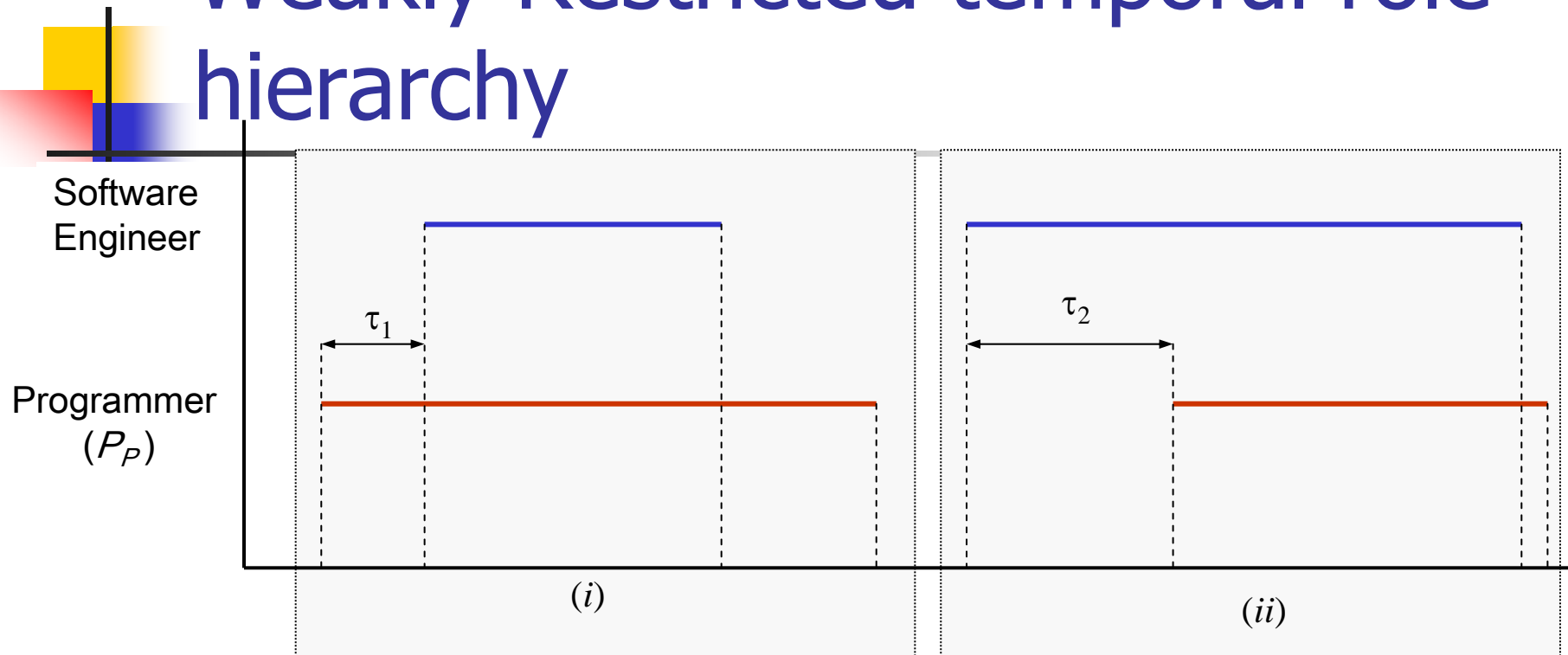
Types of role Hierarchy

- Permission-inheritance hierarchy (I-hierarchy)
 - Senior inherits juniors' permissions
 - User assigned to senior cannot activate juniors
- Role-Activation hierarchy (A-hierarchy)
 - Senior does not inherit juniors' permissions
 - User assigned to senior can activate junior
 - **Advantage:** SOD constraint can be defined on hierarchically related roles
- Activation Inheritance hierarchy (IA-hierarchy)
 - Senior inherits juniors' permissions
 - User assigned to senior can activate junior

Types of Role Hierarchy



Weakly Restricted temporal role hierarchy



Enabling intervals of Software Engineer and Programmer roles

- **Weakly Restricted:** One role needs to be enabled for the inheritance/activation semantics to apply
 - In I_{Wr} a user assigned to SE can inherit P_P in τ_2 .
 - In A_{Wr} a user assigned to SE can activate P in τ_1 .

Restricted Hierarchies

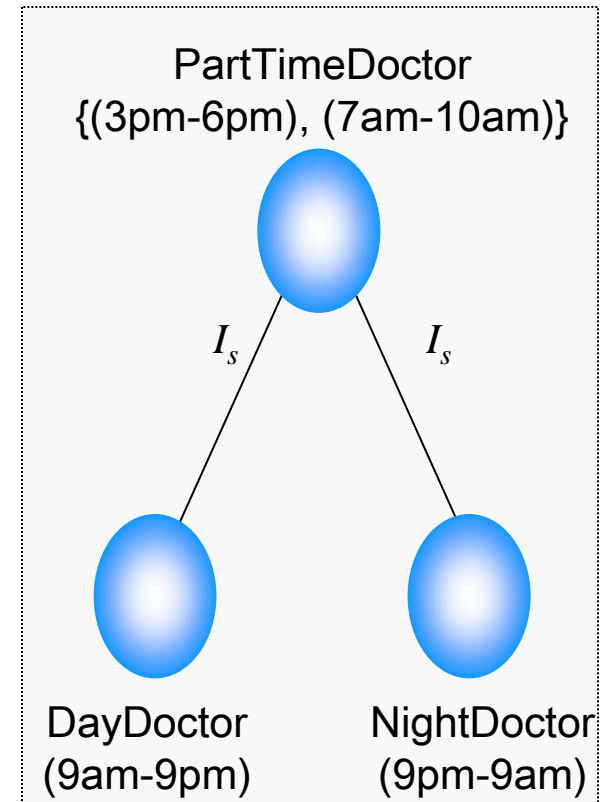
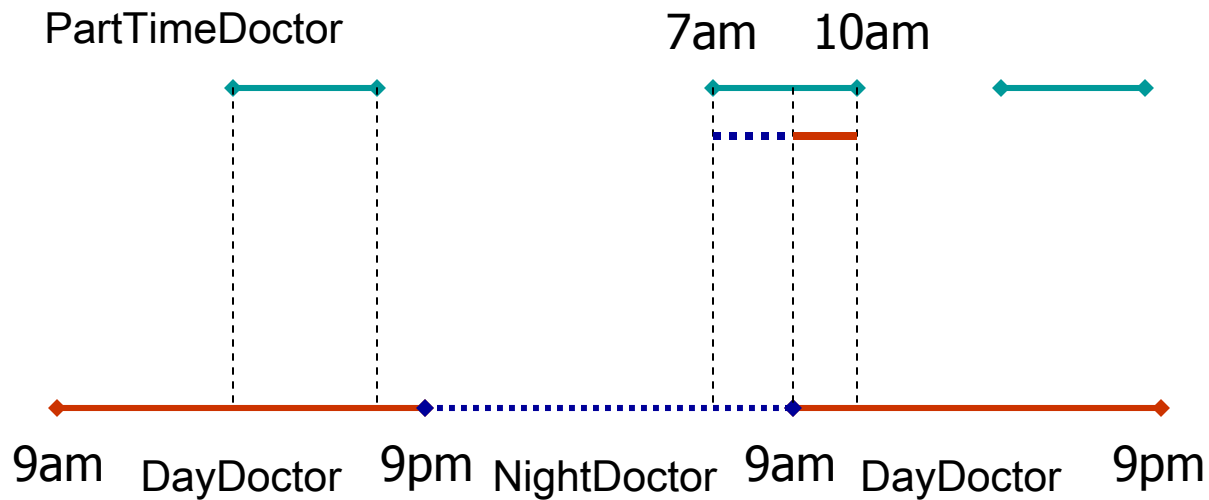
Weakly Restricted

I_w -hierarchy	$(x \geq_{w,t} y)$	$\forall p, (x \geq_{w,t} y) \wedge enabled(x, t) \wedge can_be_acquired(p, y, t) \rightarrow can_be_acquired(p, x, t)$
A_w -hierarchy	$(x \geq_{w,t}^> y)$	$\forall p, (x \geq_{w,t}^> y) \wedge enabled(y, t) \wedge can_activate(u, x, t) \rightarrow can_activate(u, y, t)$
IA_w -hierarchy	$(x \gtrsim_{w,t} y)$	$(x \gtrsim_{w,t} y) \leftrightarrow (x \geq_{w,t} y) \wedge (x \geq_{w,t}^> y)$

Strongly Restricted

I_s -hierarchy	$(x \geq_{s,t} y)$	$\forall p, (x \geq_{s,t} y) \wedge enabled(x, t) \wedge enabled(y, t) \wedge can_be_acquired(p, y, t) \rightarrow can_be_acquired(p, x, t)$
A_s -hierarchy	$(x \geq_{s,t}^> y)$	$\forall p, (x \geq_{s,t}^> y) \wedge enabled(x, t) \wedge enabled(y, t) \wedge can_activate(u, x, t) \rightarrow can_activate(u, y, t)$
IA_s -hierarchy	$(x \gtrsim_{s,t} y)$	$(x \gtrsim_{s,t} y) \leftrightarrow (x \geq_{s,t} y) \wedge (x \geq_{s,t}^> y)$

Temporal Role Hierarchy Example





Hierarchy Constraint Expressions

- Hierarchy $h \in \{I, A, IA\}$:
 - Periodicity: $(I, P, \text{enable } h)$,
 - Duration: $([I, P \mid D], D_h, \text{enable } h)$; That is $(I, P, D_h, \text{enable } h)$, $(D, D_h, \text{enable } h)$ or $(D_h, \text{enable } h)$

Example:

$(\geq_t$ is an I -hierarchy and $h = (\text{ProjManager} \geq_t \text{ProjEngineer}))$

$\text{enable } r \rightarrow \text{enable } h \text{ after } 10 \text{ min}$



Activation Constraints and Temporal Role Hierarchy

- Let P be a permission set for a Software package
 - Only 5 user licenses for the package has been obtained
 - P is assigned to **Programmer** role,
 - Suppose, we use an activation time cardinality constraint of 5 on **Programmer**
- Let **SE** be senior of **Programmer**
 - **With I or IA -hierarchy**: The use of P by more than 5 users at a time can be easily violated
 - **With A -hierarchy**: The use of P by more than 5 users at a time is controlled

Here the cardinality constraint is said to be
permission-oriented



Activation Constraints and Temporal Role Hierarchy

Requirement: At the most 5 nurses and 3 doctors on active duty; no restriction on permission use

- Apply cardinality of 5 on **Nurse** role and 3 on **Doctor** role
- Let **Doctor** be senior of **Nurse**
 - **With A-hierarchy:** 3 doctors and 5 nurses can be active at a time but, doctors will not be able to acquire **Nurse's** permissions.
 - **With I or IA-hierarchy:** 3 doctors and 5 nurses can be active at a time, and doctors will also acquire **Nurse's** permissions.

Here the cardinality constraint is said to be
user-oriented



Time-based Cardinality, Dependency and Separation of Duty Constraints

- Generic cardinality expression framework
 - Status predicates to capture all the states of GTRBAC (14)
 - Evaluation function and Projection functions
- Control flow dependency constraints
- Time-based SoD constraints
 - Systematic categorization
 - Various time-based semantics



Cardinality Constraints Examples

1	$ \Pi_1 \text{eval}(\text{enabled}(r, "t")) \geq n$	Number of roles enabled at time "t" cannot be less than n
2	$ \Pi_1 \text{eval}(\neg \text{enabled}(r, "t")) \leq n$	Number of roles disabled at time "t" cannot be more than n .
3	$ \Pi_2 \text{eval}(u_assigned("u", r, "t")) \leq n$	Number of roles assigned to "u" at time "t" cannot be more than n
4	$ \Pi_2 \text{eval}(can_activate("u", r, "t")) \leq n$	Set of roles that u can activate at time t cannot be more than n .
5	<p style="text-align: center;"> $(Daytime, \Pi_1 \text{eval}(u_assignedSet(u, "Nurse", t)) \leq n)$ Number of users assigned to Nurse role in <i>Daytime</i> cannot exceed n </p>	



Control Flow Dependency Constraints

- Typically used in workflow based systems
- Pre-condition Constraint
 - An event can happen *only if* another has already happened
 - $([I, P], \text{pre}, Y, \text{pr.}E \text{ after } \Delta t \text{ for } \Delta d)$
 - E.g., $(\text{Sat}, \text{pre}, \text{activate } \textit{Manager} \text{ for } \textit{John}, \text{enable } \textit{Employee})$
- Post-condition Constraint
 - *If* an event happens then another event *must* happen
 - $([I, P], \text{post}, Y, \text{pr.}E \text{ after } \Delta t \text{ for } \Delta d)$
 - E.g., $(\text{Sat}, \text{post}, \text{activate } \textit{SysAdmin} \text{ for } \textit{Smith}, \text{enable } \textit{SysAudit} \text{ after } 30 \text{ min})$
- Precedence
 - If two events happen then one must always precede another



GTRBAC Separation of Duty Constraints

- Important for real world commercial workflow applications and is generally used to prevent fraud
- Categorization of GTRBAC SoDs
 - Role enabling Sods
 - User-role assignment SoDs
 - Role-permission assignment SoDs
 - Activation time SoDs
 - Possibilistic role activation SoDs
 - Possibilistic permission acquisition SoDs

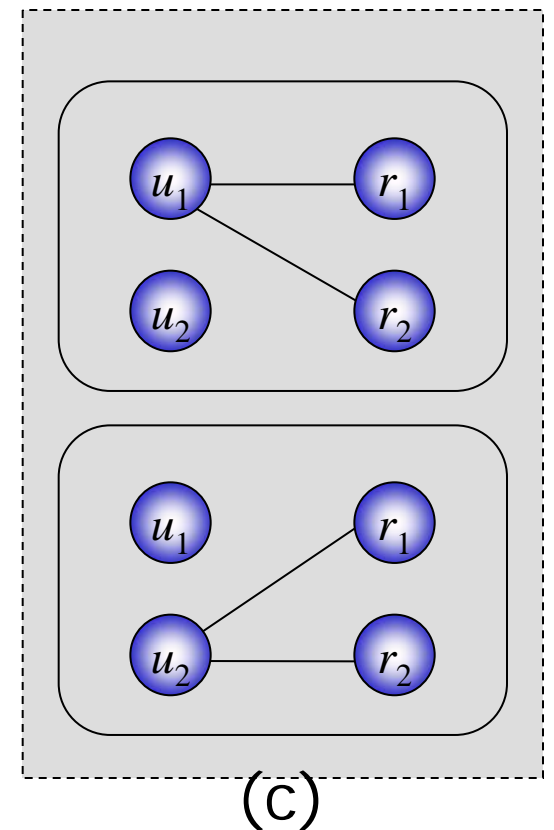
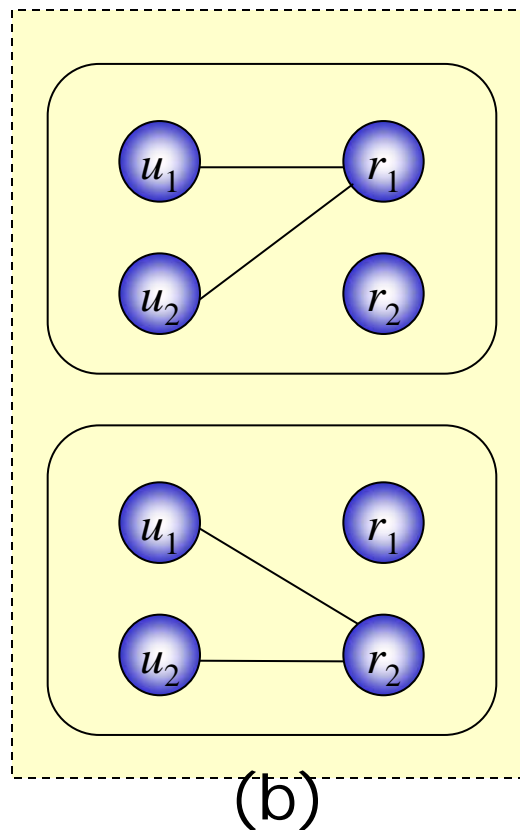
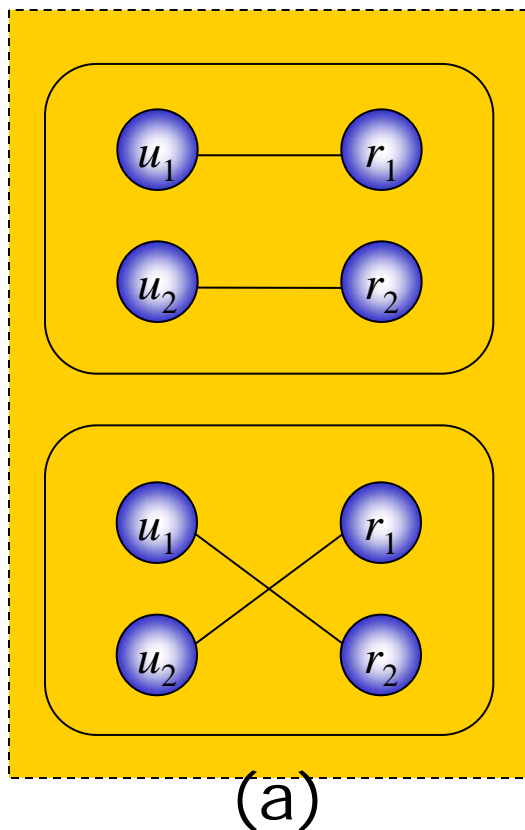


Role Enabling and User-role Assignment SoDs

- Role enabling SoD constraints (*enabled()*)
 - (I, P, EN, R) : No two roles in R can be *enabled* at the same time
 $\forall r_1, r_2 \in R, SoD \wedge enabled(r_1, t) \rightarrow \neg enabled(r_2, t)$
 - (I, P, DIS, R) : No two roles in R can be *disabled* at the same time
 $\forall r_1, r_2 \in R, SoD \wedge \neg enabled(r_1, t) \rightarrow enabled(r_2, t)$
- User assignment SoD constraints (*u_assigned*)
 - $UAS\text{-}SoD_1(I, P, UAS_1, U, R)$
 - No two roles in R can be assigned to a user in U at the same time
 - $UAS\text{-}SoD_2(I, P, UAS_2, U, R)$
 - No two users in U can be assigned a role in R at the same time
 - $UAS\text{-}SoD_3(I, P, UAS_3, U, R)$
 - Different users in U cannot be assigned different roles in R at the same time

User-role assignment SoDs

- UAS-SoD_1 does not allow c ; UAS-SoD_2 does not allow b ; UAS-SoD_3 does not allow a





User-assignment SoDs (Contd.)

- $UAS-SoD_4 = UAS-SoD_2 \wedge UAS-SoD_3$
 - Roles in R can be assigned to only one user in U
 - Example: one user must complete all the sub tasks
- $UAS-SoD_5 = UAS-SoD_1 \wedge UAS-SoD_3$
 - Users in U can be assigned only one role in R
 - Example: A team should be assigned only one consultancy job (e.g., role `ConsultantForBankA`)
- $UAS-SoD_6 = UAS-SoD_1 \wedge UAS-SoD_2$
 - A user in U can be assigned to only one role in R (and vice versa)
 - Example: A group of consultants should be assigned to different consultancy jobs (e.g., user A is assigned to role `ConsultantForBankA`, user B is assigned to role `ConsultantForBankB`, etc.)



Other GTRBAC SoDs

- Activation Time SoDs (*active()*)
 - SoDs involving active roles and sessions
 - Examples
 - No two roles in R can be in *active* state in session(s) of a user in U at the same time
 - No two users in U can have a role in R *active* at the same time
- Possibilistic Activation SoDs (*can_activate()*):
 - Captures implicit /explicit user assignments (A-hierarchy)
 - Similar to user assignment SoDs and
- Possibilistic Acquisition SoDs (*can_be_acquired()*)
 - Captures implicit/explicit permission assignment (*I*-hierarchy)
 - Example: A user in U cannot acquire different *permissions* in P at the same time.



Time-based Semantics of SoD Constraints

- Consider (I, P, UAS_1, U, R)
 - a user in U cannot be assigned to two roles in R
- (I, P, UAS_1, U, R) has various forms
 - **Weak form**: At an instant in (I, P) , if a user is assigned to a role in R , at that instant he cannot be assigned to another role in R
 - **Strong form**: For each interval in (I, P) , if there is an instant in which a user is assigned to a role, for no other instant in that interval can he be assigned to another role in R
 - **Extended Strong form**: At an instant in (I, P) , if a user is assigned to a role in R , at no other instant in (I, P) can he be assigned to another role in R



X-GTRBAC

A Policy Specification Language

- An XML conformant specification language for GTRBAC
- Allows identity or credential based dynamic assignment of roles to users
- Allows expressing multidomain policies through role mapping



X-GTRBAC

A Policy Specification Language

```
<!-- Policy Definition--> ::=  
<Policy [policy_id = "(value)"]>  
  <PolicyName> (name) </PolicyName>  
  [<!--XCredType Definition Sheet>]  
  [<!--XTemporalConstraint Definition Sheet>]  
  <!-- XML User Sheet>  
  <!-- XML Role Sheet>  
  <!-- XML Permission Sheet>  
  <!-- XML User-Role Assignment>  
  <!-- XML Role-Permssion Assignment>  
  [<!-- XSoD Definition Sheet>]  
  [<!-- XHierachy Definition Sheet>]  
  [<!-- Local Policy Definitions-->]  
  [<!-- Policy Relationship Definitions>]  
</Policy>
```



An XML instance of XUS

```
<XUS>
  <Users>
    <User user_id="j1">
      <UserName >John</ UserName >
      <CredType cred_type_id ="C100">
        < type_name >Nurse</type_name>
        <CredExpr>
          <age> 30 </age>
          <field> ophthalmology </field>
          <level> 5 </level>
          <status> single </status>
        </CredExpr>
      </CredType>
      < MaxRoles>2</MaxRoles>
    </User >
    ....
  </Users>
</XUS>
```



An XML instance of XRS

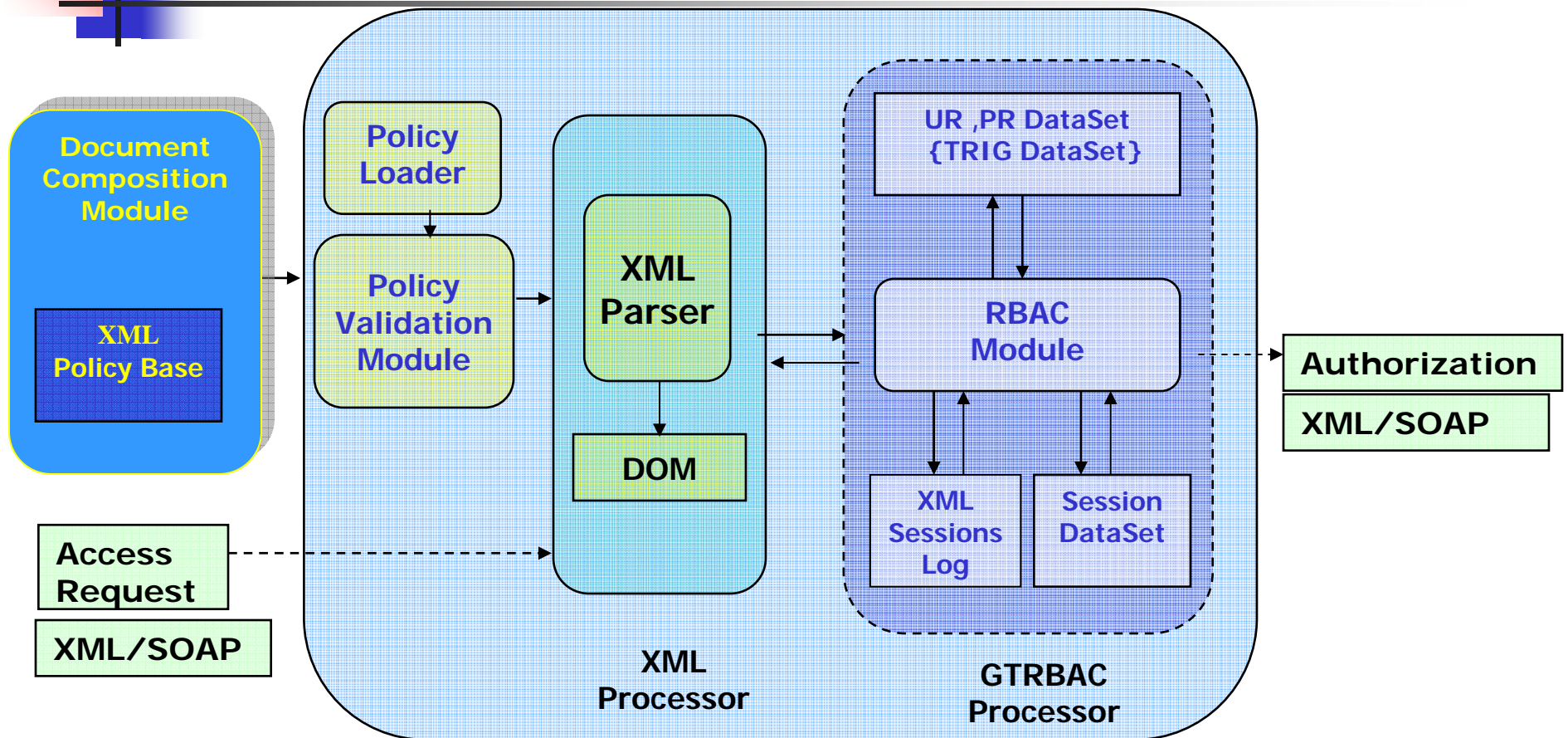
```
<XRS>
  <Role role_id = "R100">
    <RoleName> Nurse </RoleName >
    <Senior HType = "IA"> Eye_Doctor
    </Senior>
    <Cardinality> 8 </Cardinality >
  </Role>
  <Role role_id = "R200">
    <RoleName> Eye_Doctor </RoleName>
    <Junior HType = "IA"> Nurse
    </Junior>
    <Senior HType = "A"> Eye_Surgeon
    </Senior>
    <Cardinality> 6 </Cardinality>
  </Role>
</XRS >
```


Periodicity and Duration Expressions

```
<XTempConstDef>
  <PeriodicTimeExpr pt_expr_id="PTQuarterWeekOne"
    i_expr_id="Year2003">
    <StartTimeExpr>
      <Year>all</Year>
      <MonthSet>
        <Month>1</Month>
        <Month>4</Month>
        <Month>7</Month>
        <Month>10</Month>
      </MonthSet>
      <WeekSet>
        <Week>1</Week>
      </WeekSet>
    </StartTimeExpr>
  </PeriodicTimeExpr>
  .....
</XTempConstDef>
```

```
<DurationExpr
  d_expr_id="SixWeeks">
  <cal>Weeks</cal>
  <len>6</len>
</DurationExpr>
```

X-GTRBAC Architecture



Validation support is provided by Apache Xalan XSLT engine built into JAXP



Policy Design Issue

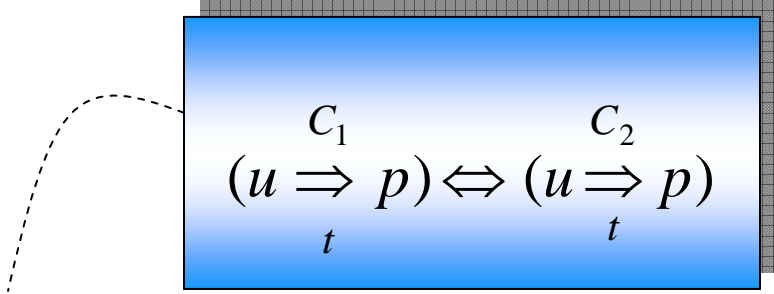
- GTRBAC constraint set is not minimal
- Constraint design considerations
 - Usability: Clarity of Semantics, Manageability
 - Complexity of specification

Complexity parameter	Description
$n.R$	n roles
$n.S$	n default assignments
$n.T_r$	n temporal constraints on (n) roles
$n.T_{ur}$ ($n.T_{rp}$)	n temporal constraints on user- assignment
$n.A_{ur}$ ($n.A_r$)	n per-user-role (per-role) time constraint
$n.H$	$n.H$ indicates n hierarchical relations

Activity-equivalence

- GTRBAC Configuration C :
 - $(T, \text{Users, Roles, Permissions, RH})$, where T is the set of constraints, RH is a A -hierarchy.
- Activity-equivalent configurations

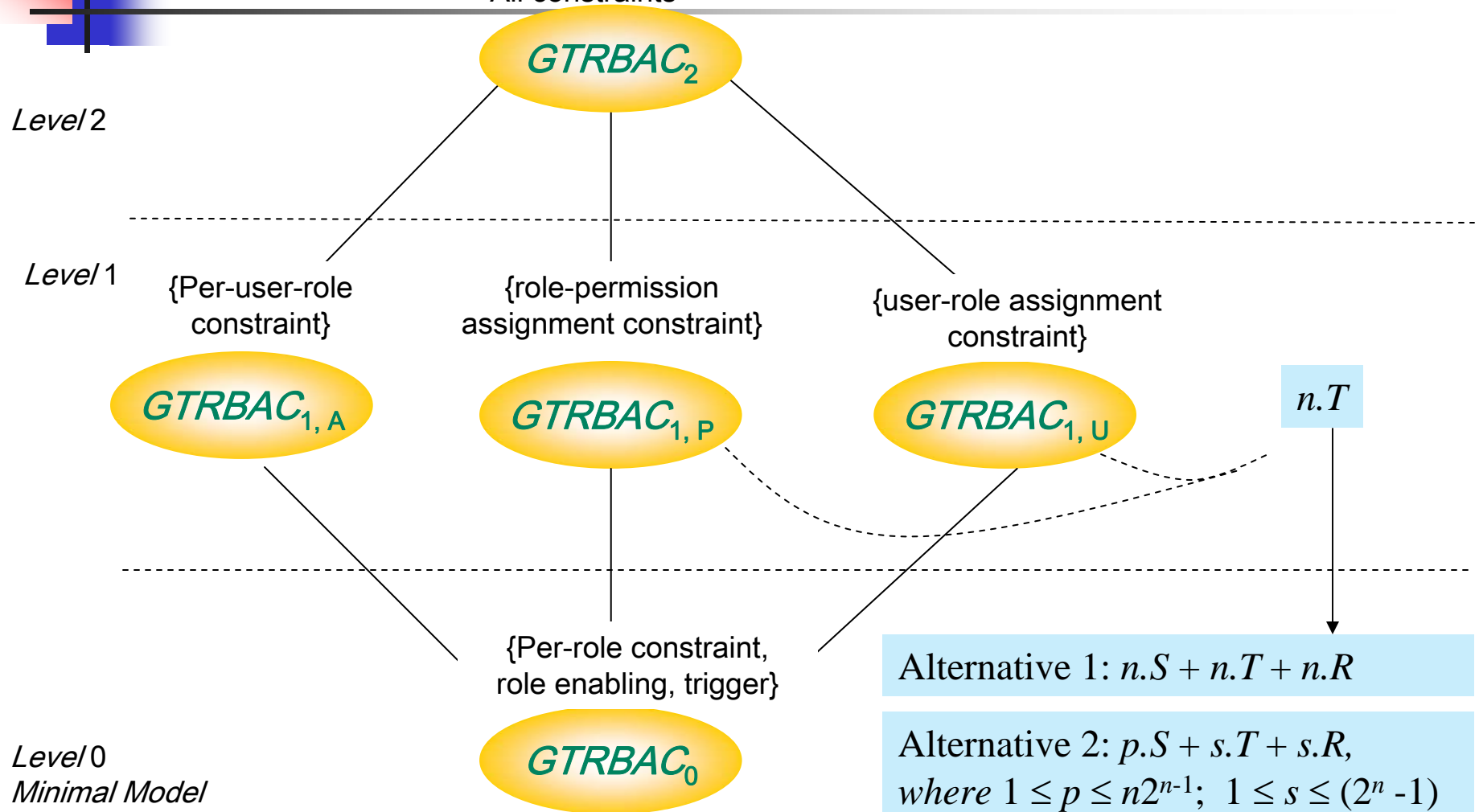
- $C_1 \approx C_2$:


$$(u \underset{t}{\Rightarrow} p) \Leftrightarrow (u \underset{t}{\Rightarrow} p)$$

- u can acquire p at t in C_1
- There exists a minimum set of constraint types

GTRBAC Family of Models

All constraints



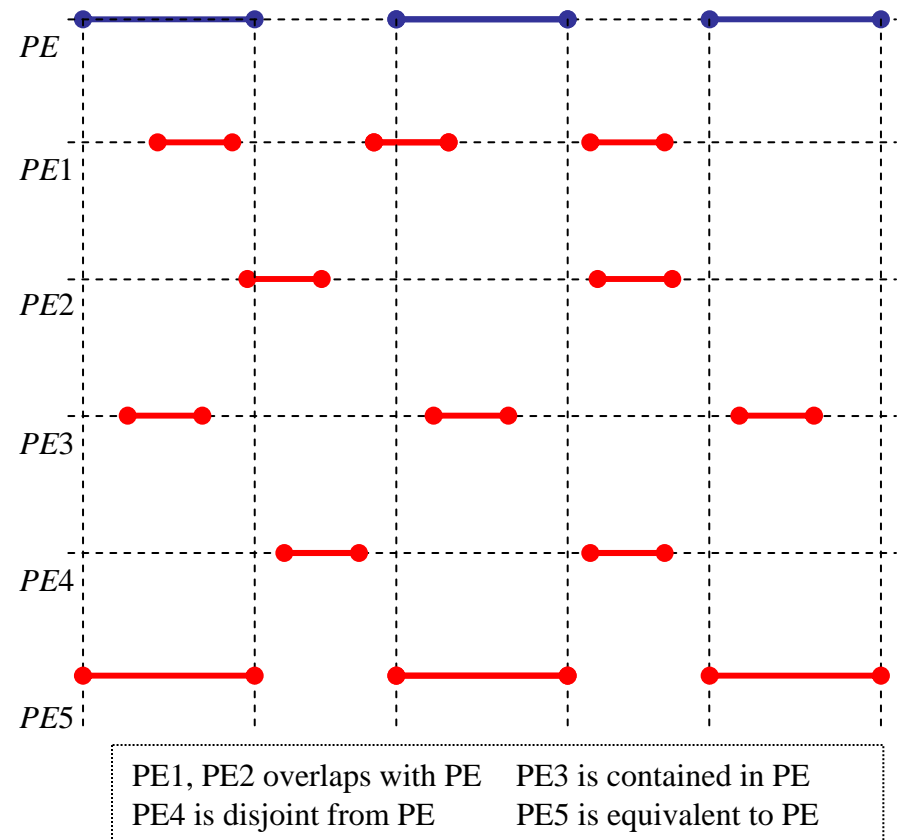


Approach to transformation

- Temporal constraint on user-role assignment is changed to role enabling constraint
- Two approaches
 - Simple approach:
 - PE_u of each $u-r$ assignment becomes enabling constraint for r
 - Minimal Disjoint Set Approach:
 - PE_u of each $u-r$ assignment becomes enabling time constraint for a set of new roles
 - First compute Minimal Disjoint Set (MDS) of all PE_i s
 - Create new roles associated with each element of MDS
 - User u is assigned to new roles that have elements of MDS that correspond to PE_u
- Same for the permission role assignment

MDS Approach: Relations on periodic expressions

- **Containment:** PE1 is contained in PE2
 - All instants in PE1 is in PE2
- **Equivalence:** PE1 and PE2 have same time instants
- **Overlapping:** At least one time instant is common to both PE1 and PE2
- **Disjoint:** No common time instants in PE1 and PE2





Minimal Disjoint Set (MDS) & Minimal Subset (MS) of PE over MDS

- Let $PE = \{PE_1, PE_2, \dots, PE_n\}$
 - $MDS_{PE} = \min_m \{PE'_i \mid 1 \leq i \leq m\}$ such that
 - $PE'_1, PE'_2, \dots, PE'_m$ are pair-wise disjoint
 - $PE'_1 \cup PE'_2 \cup \dots \cup PE'_m = PE_1 \cup PE_2 \cup \dots \cup PE_n$,
 - If PE'_i contains a time instant of PE_j then it does not contain a time instant that is not in PE_j
 - Bounds: $1 \leq |MDS_{PE}| \leq (2^n - 1)$
- MS of $PE_j \in PE$ over the MDS_{PE}
 - subset of MDS_{PE} that collectively contains all the time instants of PE_j
 - Bounds: $n \leq \sum |MS_{PE_j}| \leq n2^{n-1}$



Example

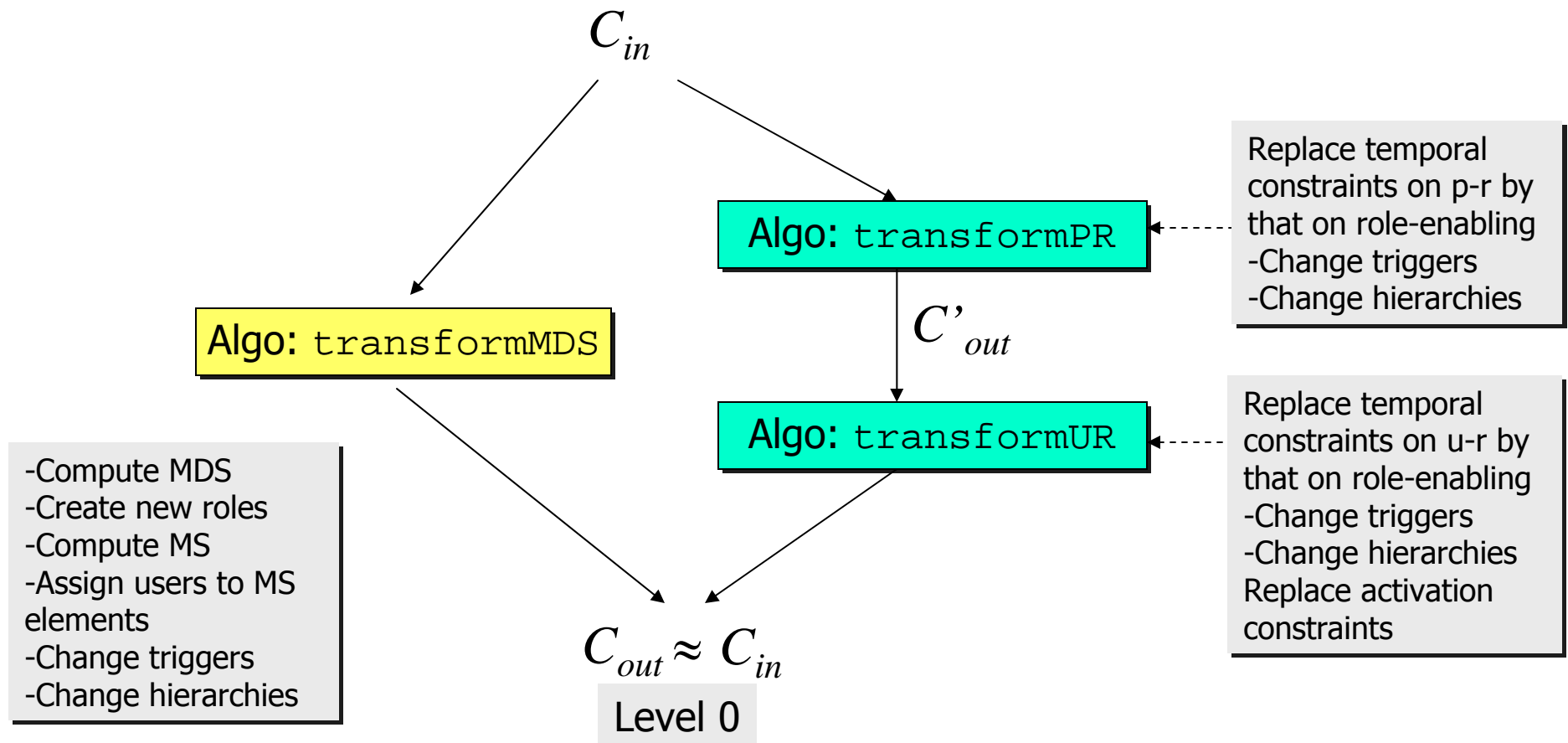
- Let $PE = \{PE_A, PE_B, PE_C, PE_D, PE_E\}$, where
 - $PE_A = \{\text{Sun, Mon, Tue, Wed, Thu, Fri}\}$;
 - $PE_B = \{\text{Sun, Tue}\}$,
 - $PE_C = \{\text{Sun, Tue, Thu, Fri}\}$,
 - $PE_D = \{\text{Sun, Mon, Tue, Wed, Sat}\}$,
 - $PE_E = \{\text{Thu, Fri}\}$.
- *MDS of PE* is
 - $\{PE_1, PE_2, PE_3, PE_4\} = \{\{\text{Sun, Tue}\}, \{\text{Thu, Fri}\}, \{\text{Mon, Wed}\}, \{\text{Sat}\}\}$
- MS values are as follows:
 - *MS of PE_A* : $\{PE_1, PE_2, PE_3\}$;
 - *MS of PE_B* : $\{PE_1\}$
 - *MS of PE_C* : $\{PE_1, PE_2\}$;
 - *PE_D* : $\{PE_1, PE_3, PE_4\}$
 - *PE_E* : $\{PE_2\}$

Algorithm TransformMDS – Replacing temporal constraint on user-role assignment by temporal roles

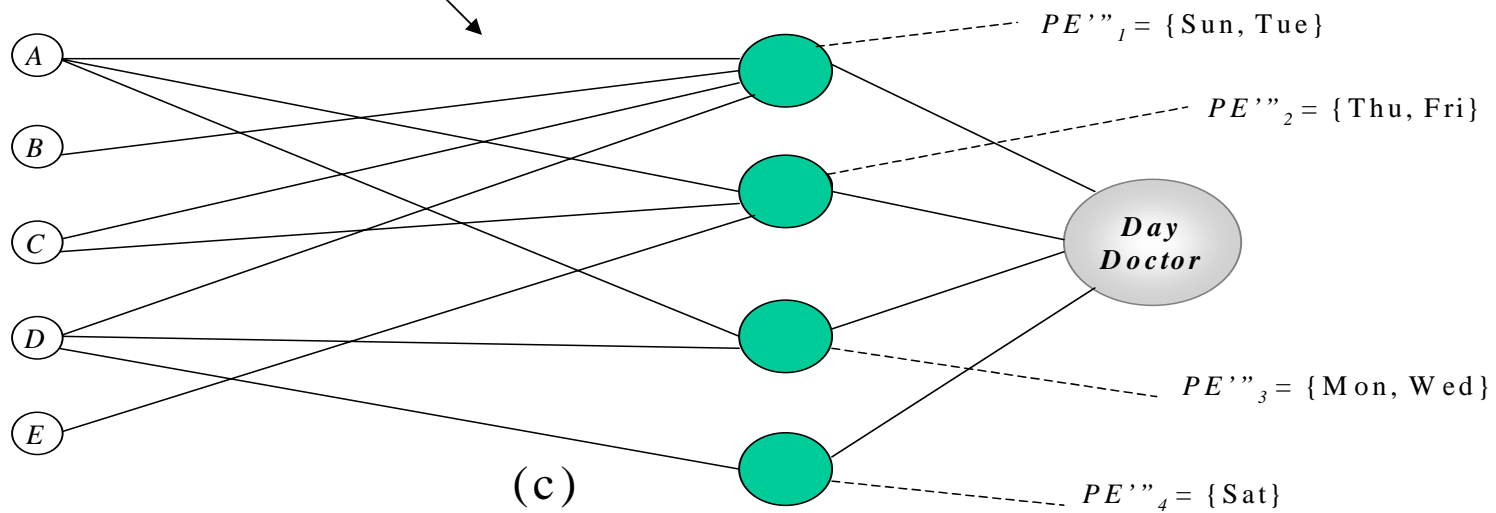
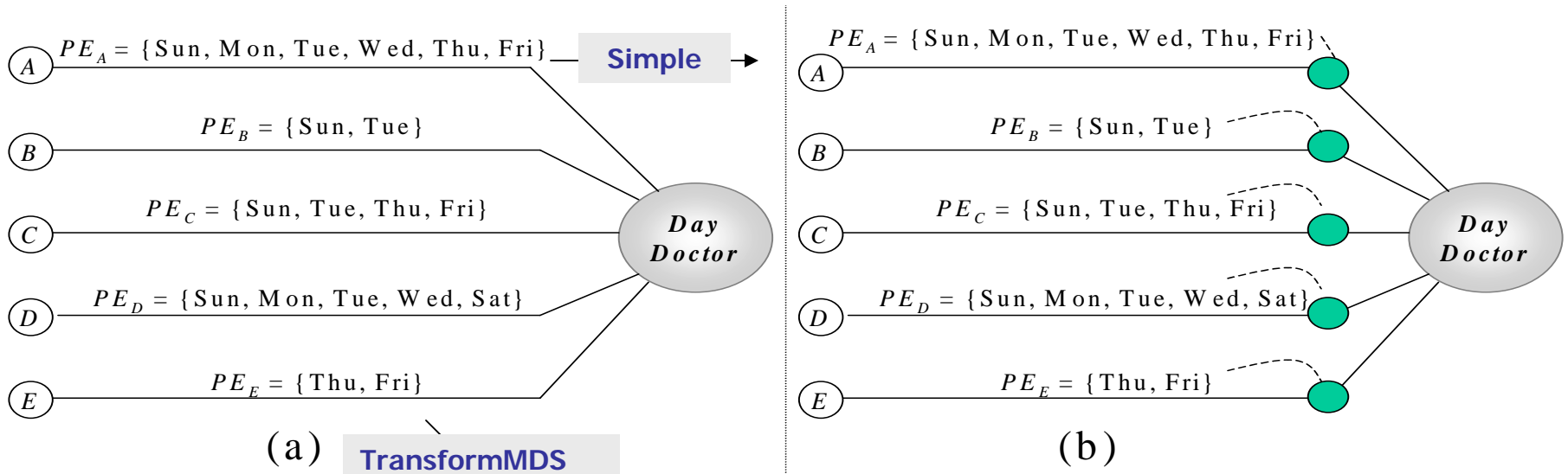
Input: C_{in} **Output:** C_{out}

1. $C_{out} = \{T', Roles', RH'\} = C_{in} = \{T, Roles, RH\};$
2. **FOR** each $r \in Roles$ **DO**
3. Let $PE = \{PE_1, PE_2, \dots, PE_n\}$ & $U = \{u_1, u_2, \dots, u_n\}$ be s.t. $(PE_i, \text{assign } r \text{ to } u_i) \in T'$;
4. Compute *MDS* of PE ; Let the computed $MDS = \{PE'_1, PE'_2, \dots, PE'_n\};$
5. **FOR** $i = 1$ to n **DO**
6. Compute MS_{PE_i} for PE_i
8. **FOR** each $PE'_i \in MDS$ **DO**
9. Create a unique role r_i ;
10. **FOR** all $u_k \in U$ such that $PE'_i \in MS_{PE_k}$ for PE_k **DO**
11. Add *default* assignment (assign r_i to u_k) in T' .
12. Add constraint $(PE_i, \text{enable } r_i)$ in T' .
13. Remove constraint $(PE_i, \text{assign } r \text{ to } u_i)$ from T' ;
14. $Roles' = Roles' \cup \{r_i\};$
15. $RH' = RH' \cup \{r < r_i\};$

Simple & MDS approaches



Example



Alternative (c): $p.S + s.T + s.R + s.H$,
 where $n \leq p \leq n2^{n-1}$; $1 \leq s \leq (2^n - 1)$

Alternative (b): $n.S + n.T + n.R + n.H$



Design Guidelines

- $\text{GTRBAC}_{1,U}$ is better than *alternative b* as the policy representation is less complex in terms of number of roles and hierarchies
- $\text{GTRBAC}_{1,U}$ is flexible – one can schedule role enabling and user assignments separately
- when p_n and s_n are close to n and 1, *alternative c* may be better than $\text{GTRBAC}_{1,U}$ representation

Alternative (b):

$$n.S + n.T_R + n.R + n.H$$

Alternative (c):

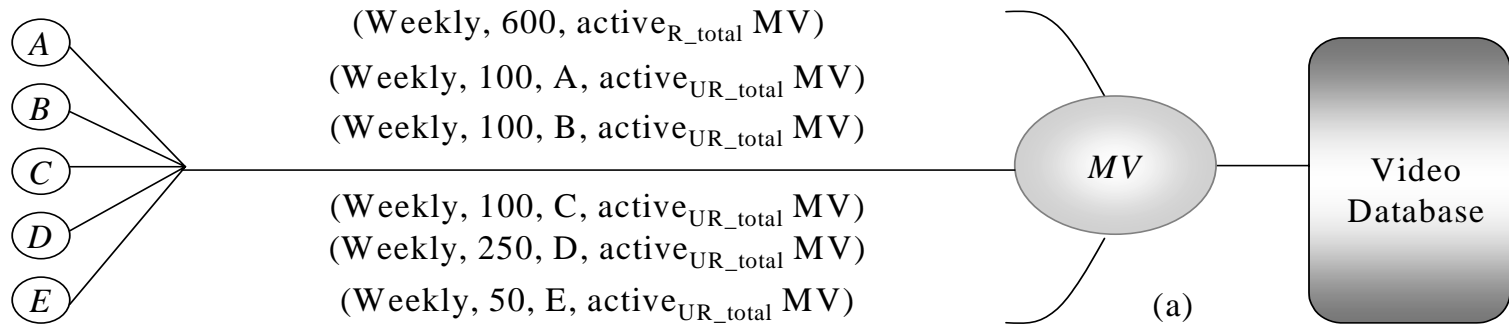
$$p.S + s.T_R + s.R + s.H$$

where

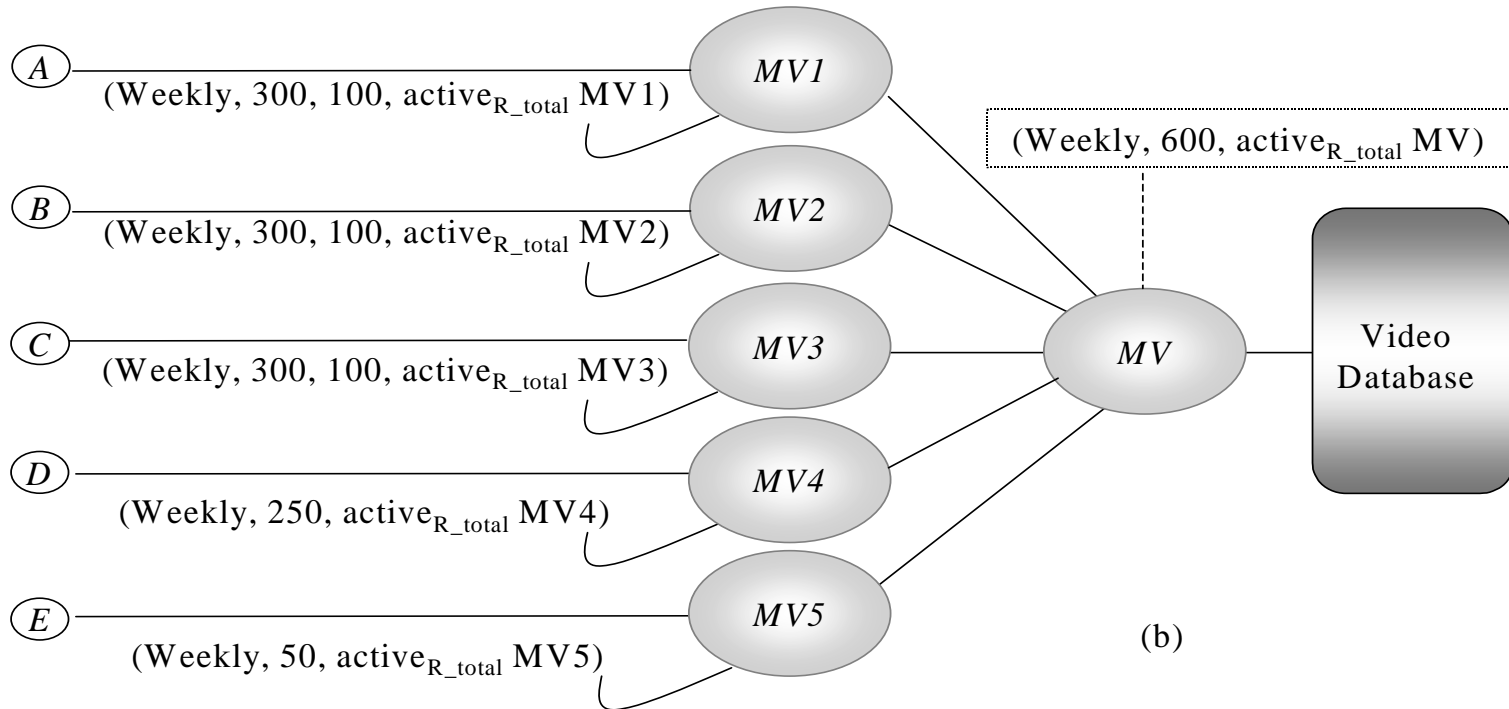
$$n \leq p_n \leq n2^{n-1};$$

$$1 \leq s_n \leq (2^n - 1)$$

Replacing *per-user-role* by *per-role*

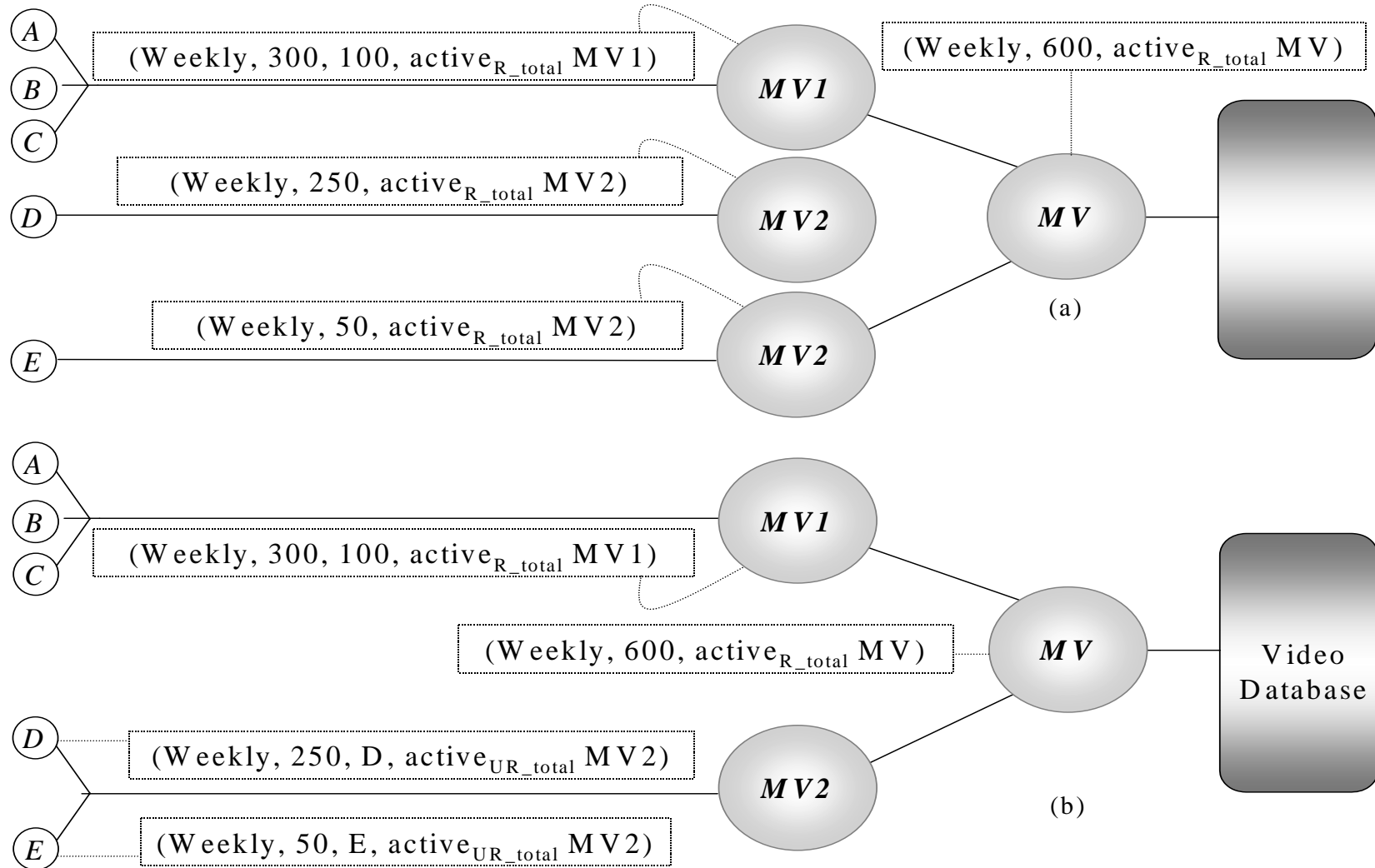


Alternative (a): $n.A_{ur}$, where A_{ur} user-role activation



Alternative (a): $n.A_r$, $n.R + n.H$

Replacing *per-user-role* by *per-role*



$GTRBAC_{1,A}$ representation: $(n_x - n_y).A_{UR} + n_y.A_R + c.(b.n_y + 1).(R + H)$;

where, (1) $n_x = |D_m|$ and $n_y = |D'|$, such that (1) $D' \subseteq D_m$ and (2) if $d \in D'$ then $C_m(d) > 1$; (2) $b = 1$ if $(n > n_x)$; $b = 0$ otherwise; (3) $c = 1$ if $(n > n_x > 0)$; $c = 0$ otherwise.



Design Guidelines

- Per-role constraint with default value
 - If there are many common durations,
- per-user-role constraint
 - Per-user requirements vary significantly
 - More flexibility (e.g., requirements vary constantly)
- Hybrid approach (b in previous slide) can give balanced representations



Related Work

- OASIS RBAC Model [Bacon, 02]
 - Precondition on role activation to support active security
 - No triggers
- RSL2000 Constraint Specification Language [Ahn et. al., 2000],
- Need for activation hierarchy [Sandhu, 1999]
 - Identified its usefulness in expressing MAC
- Separation Duty Constraints
 - Listing of useful set of SoD Constraints
 - [Simon et. al., 1997],[Gligor et. al., 1998]
- None address timing issues
- GTRBAC SoDs subsume all the SoD constraints identified
- GTRBAC Triggers and SoDs provide a technical foundation for enforcing history based SoD constraints



Conclusion

- Role based access control can be used to support diverse set of access control requirements
- Time-based access is a crucial requirement in emerging applications
- GTRBAC's constraint set is not minimal—however, they are beneficial for practical use



Current and Future work

- Secure Interoperation
 - Integer Programming approach (for tightly coupled environments)
 - Trust-based access management (loosely coupled environment)
 - Issues related to Grid, P2P
- GTRBAC - extended to LoT-RBAC
 - Implementation in mobile environment (near completion)
- Policy evolution and hybrid hierarchy management
 - Administrative tools and techniques
- Extension of the policy design work to generate tools for efficient RBAC policy administration



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