## IS 2610: D ata Structures

Elementary Data Structures

Jan 12, 2004

D ata Type

- Is a set of values and a collection of operations on those values.
- Inbuilt data types
- Int
- Float
- Character
- New Data types
- Define values to operate (arguments of a function)
- Define operation (function definition)


## Sample function definition

```
#include <stdio.h>
int lg(int);
main() {
    int i, N;
    for (i = 1, N = 10; i <= 6; i++, N *= 10)
        printf("%7d %2d %9d\n, N, lg(i), N*lg(N))
}
Int lg(int N){
    int i;
    for (i = 0;N > 0; i++, N/= 2);
    return i;
}
```


## D ata Structure

- Goal is to build data structures that allow us to handle collections of data
- What operations need to be performed?
- How to implement these operations?
- Simplest way to organize data in C
- Arrays
- Structures


## Software Engineering practice

- Interface (header file)
- Defines data structures
- Declare functions to be used to manipulate the data structure
- Implementation (separate c file)
- Of the functions declared in Interface
- Client (main application program)
- Program that uses the functions declared in the Interface to work at a higher level of abstraction


## Arrays

- Most fundamental data structure
- Fixed collection of same-type data
- Access is made by using an index
- Contiguously stored
- Direct correspondence with memory systems
- Entire memory can be considered as an array of memory locations, with memory addresses corresponding to the array indices
- In C array definition
- int A1[N]; int A2[N][M]; char str[50];
- A1[4]? A1[i] = *(A1+i)?
- Suppose you have to pass huge array as an argument?


## Array

- Dynamic Memory Allocation

```
#define N 1000
main() {
    int i, a[N];
#nclude <stdlib.h>
main(int argc, char* argv) {
    int i, N = atoi(argv[1]);
    int *a = malloc(N*sizeof(int);
    if (a==NULL) Insufficient memory
}
```

- Sieve of Eratosthenes

```
#define N 20
main() {
    int i, j, a[N];
    for (i = 1; i<N; i++) a[i]=1;
    for (i = 2; i<N; i++) Finding primes
        if (a[i])
        for (j = i; i*j<N; j++) a[i*j] = 0;
    for (j = 2; j<N; j++)
        if (a[i]) printf ("%4d \n", i);
}
```


## Linked List

- A set of items where each item is part of a node that also contains a link to a node
- Self referent structures
- Cyclic structures possible
- C code
typedef struct node *link;
struct node \{char ch; link next; \}
link $h=$ malloc(sizeof $* h$ );
note that h->next denotes the $2^{\text {nd }}$ node and $h->c h$ denotes The value "a"
h



## List traversal

- Print each element of the list
h

- Use while loop
- Use for loop
- Inverting a list?


## Linked Lists



- Delete Operation
x (delete after x )
h

- Exchange Operation
h



## D oubly Linked List

```
typedef struct node *link;
struct node {char ch; link prev; link next;}
link h = malloc(sizeof *h);
```


(delete after h)

## String

- Variable length array of characters
- Has a starting point and a string-termination character ('\0') at the end
- Array based implementation in C
- Array of characters different from string - associated with length
- String length may change
- Many applications involve processing textual data
- Computers provide access to bytes of memory that correspond directly to characters of strings


## Common String functions

- strlen(a)
for (i=0; a[i] != 0; i++); return $i$;
- $\operatorname{strcpy}(\mathrm{a}, \mathrm{b})$
for $(i=0 ;(a[i]=b[i])!=0 ; i++) ; \quad$ while $(* a++=* b++) ;$
- $\operatorname{strcmp}(a, b)$
for $(i=0 ;(a[i]==b[i])!=0 ; i++)$;
if $(a[i]=0)$ return 0 ;
return $\mathrm{a}[\mathrm{i}]-\mathrm{b}[\mathrm{i}]$
- $\operatorname{strcat}(a, b)$
$\operatorname{strcpy}(a+\operatorname{strlen}(a), b)$


## Abstraction

- Layers of abstraction
- Abstract model of a bit with binary 0-1 values
- Abstract model of a machine from from dynamic properties of the values of a certain set of bits
- Abstract model of a programming language that we realize by controlling the machine with a machine -language program
- Abstract notion of algorithm implemented in C
- Abstract data types
- Develop abstract mechanisms for certain computational tasks at a higher level
- New layer of abstraction
- Define objects we want to manipulate
- Represent data in data structures
- Define operations that we perform on them
- Implement the algorithm


## Abstract D ata Type

- A data type that is access only through an interface
- Refer to a program that uses ADT as a client and program that specifies the data type as an implementation
- Interface is opaque - clients cannot see implementation
- Benefits of ADTs
- Provide an effective mechanism for organizing large software systems
- Provide ways to limit the size and complexity of interface between algorithms and associated data structures and programs that use the algorithms and data structures
- ADTs interface defines a precise means of communication


## Pushdown Stack ADT

- An ADT that comprises two basic operations: insert (push) a new item, and delete (pop) the item that was most recently inserted
- Last in- first out (LIFO Queue)



## Pushdown-stack ADT interfaces

- Use in evaluation of arithmetic expression - Infix expression (customary way)
- Operator comes between the operands
- $4+5$ is written as $45+$
- Postfix expression
- Operator comes after the operands
- $4+5$ is written as $45+$
- Postfix expression
- Operator comes after the operands
- $4+5$ is written as $45+$
void STACKinit(int); int STACKempty(); void STACKpush(Item); Item STACKpop();
- Interfaces: Client may use the four operations - store in STACK.h


## Postfix notation

- What is the postfix for the following infix - $6+5$ * 9 ?
- What is the infix for the following postfix
- $598+46^{* *} 7$ + *?
-598-71-*+7*?
- Note parentheses are not necessary in postfix


## Postfix notation and Pushdown Stack



## Stack Implementation (Array)

```
static Item *s;
static int N;
void STACKinit(int maxN)
    {s = malloc(maxN*sizeof(Item));N = 0;}
    int STACKempty()
    {return N==0;}
void STACKpush(Item item)
    {s[N++] = item;}
Item STACKpop()
    { return s[--N];}
```


## Stack Implementation (Linked-list)

- Assume auxiliary function

```
typedef struct STACKnode* link;
struct STACKnode {Item item; link next;}
static link head;
link NEW(Item item, link next;}
    { link x = malloc(sizeof *x);
        x->item = item; x->next = next;
        return x;
}
```

- Write the functions


## First-In First O ut Q ueues

- An ADT that comprises two basic operations: insert (put) a new item, and delete (get) the item that was least recently used

```
typedef struct QUEUEnode* link;
struct QUEUEnode {Item item; link next;}
static link head;
link NEW(Item item, link next;}
    { link x = malloc(sizeof *x);
        x->item = item; x->next = next;
        return x;
    }
```


## First-class ADT

- Clients use a single instance of STACK or QUEUE
- Only one object in a given program
- Could not declare variables or use it as an argument
- A first-class data type is one for which we can have potentially many different instances, and which can assign to variables whichcan declare to hold the instances


## First-class data type - Complex numbers

- Complex numbers contains two parts
- $(a+b i)$ where $i^{2}=-1$;
a $(a+b i)(c+d i)=(a c-b d)+(a d+b c) i$

Typedef struct \{float r; float i;\} Complex;
Complex COMPLEXinit(float, float)
float Re(float, float);
float Im(float, float);
Complex COMPLEXmult(Complex, Complex)

```
Complex t, x, tx;
    t= COMPLEXInit(cos(r), sin(r))
    x = COMPLEXInit(?, ?)
    tx = COMPLEXInit(t, x)
```


## First-class data type - Q ueues

```
typedef struct queue *Q;
void QUEUEdump(Q);
        Q QUEUEinit(int);
        int QUEUEempty(Q);
void QUEUEput(Q, Item);
Item QUEUEget(Q);
```

Q queues[M];
for (i=0; i<M; i++)
queues[i] = QUEUEinit(N);
printf("\%3d ", QUEUEget(queues[i]));

## Recursion and Trees

- Recursive algorithm is one that solves a problem by solving one or more smaller instances of the same problem
- Recursive function calls itself
- Factorial?
- Fibonacci numbers?

Euclid's method for finding the greatest Common divisor
int gcd(intm, int n)\{
if ( $\mathrm{n}==0$ ) return m ; return $\operatorname{gcd}(\mathrm{n}, \mathrm{m} \% \mathrm{n})$;

## Tree traversal (binary tree)

- Preorder
- Visit a node,
- Visit left subtree,
- Visit right subtree
- Inorder
- Visit left subtree,
- Visit a node,
- Visit right subtree
- Postorder
- Visit left subtree,

- Visit right subtree
- Visit a node

