

Threat Analysis in Online Social Network Systems

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Purpose

- Take online social network systems as examples to demonstrate how to conduct a threat analysis in a complex system
- Investigate & analyze various security & privacy issues in the most popular online social network systems, such as Facebook, LinkedIn, Foursquare and Yelp
- Be aware of these problems & know how to mitigate or avoid the potential attacks



Steps of Threat Analysis





Coverage

- Not focus on the traditional problems
 - Authentication
 - Secure Communication
 - Web-based Attacks; E.g., SQL Injection, Cross Site Scripting
- Focus is on the new vulnerabilities that exist in online social networks
 - Traditional online social networks (OSN); E.g., Facebook & LinkedIn
 - Location-based social networks (LBSN); E.g., Foursquare & Yelp



Steps of Threat Analysis







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OSN

Friendship/Social Network



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Entities, Elements & Mechanisms

- User's Social Network
 - Friends
 - Mutual Friends
 - Recommended Friends
- User's Posts
 - Messages
 - Photos
 - Check-ins (LBSN)

- User Identity / User Profile
 - Attributes
- Venue (LBSN)
 - Attributes

- Mechanisms
 - User Authentication
 - Access Control Mechanisms



Steps of Threat Analysis





Investigation of User Identity / User Profiles & Venues





Cloned Identity





Identity Clone Attack [6] - Design

- Attributes: name, education, birthday...
- Friend network
 - Friend List (FL): Connected friends of an ID
 - Recommended Friend List (RFL):
 - ✓ Generated by OSN systems (function of "*People You May Know*" on Facebook)
 - \checkmark Share same RFs
 - Excluded Friend List (EFL):
 - \checkmark Social embarrassments
 - \checkmark Attackers try to connect these individuals



What are the best targets

Not having Account

Inactive Account

Popular / Authority Account

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Friend Networks As Target





Cloned Identity Detection



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Profile Similarity

Attribute Similarity

 $S_{att}(P_{c}, P_{v}) = \frac{SA_{cv}}{\sqrt{|A_{c}| \times |A_{v}|}}$

Basic Principle: Similar Attributes in Two Profiles

Friend Network Similarity For Basic Profile Similarity (BPS)

$$S_{bfn}(P_{c}, P_{v}) = (\alpha S_{ff} + \beta S_{frf} + \gamma S_{fef})$$

Basic Principle: Mutual Friends in Friend Networks

For Multiple-faked Identities Profile Similarity (MFIPS)

 $S_{mfn}(P_{c},P_{v}) = \alpha(S_{s-ff} + S_{s-cf}) + \beta(S_{s-frf} + S_{s-cfrf}) + \gamma S_{s-fef}$

Basic Principle: Similar Friends in Friend Networks



Experiments



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Infer User's Profile Information

- Assumptions: Friends tend to share the same interests
- Inferring a targeted user's private attribute based on his/her friends' public attributes
- Example:
 - A user hides his education and occupation from the public
 - Many of a user's friends are current students at the University of Pittsburgh
 - Inference: University of Pittsburgh, Student



Venue Attacks in LBSNs [2]

- Venue Attributes
 - Creator
 - Owner
 - Name
 - Address
 - Geo-location
 - Category
 - Statistical Information Owner
 - Promotion/Coupon (Set by Owner)





Malicious Venue Creation Attack

- ANY user can create ANY type of a venue without being subjected to any AUTHENTICATION and the AUTHORIZATION from the actual owner
- Venue Not Created in a LBSN
 - Does not exist in the real world: deceive and confuse users, destroy users' trust for LBSNs
 - Exists in the real world but not willing to share; e.g. home, private place
- Venue Already Created in a LBSN
 - Create a similar venue using a similar/alternative name; e.g., School of Information Sciences - iSchool



Venue Ownership Hijacking Attack

- Bypass the owner authentication process & become the owner of the created venue
- Owner Authentication in Foursquare, Yelp and Facebook Place
 - Phone number
 - Address
- Impacts

. . .

- Expose customers' visit information: users' privacy
- Manipulate coupons/promotions: financial loss and/or destroy user trust on the venue
- Change the address of the venue





Venue Location Hijacking Attack

- Venue's location is associated with its geo-location not the physical address
- Geo-location is dynamic in terms of possible inaccurate GPS signals
- Location update: the center of all the honest check-ins marked by a LBSN







Users' Honest Check-ins & Marked as Host Check-ins by System



Users' Dishonest Check-ins & Marked as Honest Check-ins by System



Actual Location of the Venue



Users' honest Check-ins & Marked as Dishonest Check-ins by System



Users' Dishonest Check-ins & Marked as Dishonest Check-ins by System



Manipulated Location of the Venue



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The Movements of the Locations of the LERSAIS Lab



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Combined Venue Attacks







Moved 2 Miles away in May, 2012

Moved 3 Miles away in July, 2012



New Venue Created & Its Check-ins in August, 2012



Investigation of User's Social Network & Posts



- 1. Mutual-friend Based Attack
- 2. Friend Inference Attack

Resource Sharing

Issues



Issues Related to Users' Friend Lists

- Importance of the friend list
- What a user's friends reveals
 - Family, Work, Income, Reputation, Religion...
 - Used for Identity Clone Attacks
 - Used for Inferring Private Attributes





Attacks - Expose a User's Social Network

- Mutual-friend based Attack
- Friendship Identification and Inference Attack





Mutual Friend Feature

- Show mutual friends between two users
- Useful feature, *e.g.* Friend Recommendation, Friend Introduction

Lack of the Access Control Mechanism !





Attack Example





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Definition of a Mutual-friend based Attack

- Attacker (a)
- Target (<u>*t*</u>):
 - ✤ <u>t</u> has privacy settings for <u>a</u>
 - * <u>a</u> does not know <u>t</u>'s friends and distant neighbors
- **Knowledge of an Attacker (Assumptions):**
 - ✤ <u>a</u> knows his friend list
 - ✤ a can find t
 - For each \underline{u} that \underline{a} can find, \underline{a} can query $MF(\underline{a}, \underline{u})$
- **Mutual-friend based Attack:** •
 - * at least one of <u>t</u>'s friends and/or distant neighbors are **exposed** to <u>a</u> by querying mutual friends





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Attack Structures

- Mutual-friend based Attack
 - Conduct various well-designed mutual friend queries
- Queries are designed based on attack structures
 - Types of Attack Structures
 - → Used to identify a target's friends (*BASFs*)
 - ➢ Used to identify the target's distant neighbors (*BASDNs*)



BASFs



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Specific-Target Attack

- Attacker (\underline{a}) has a specific target (\underline{t})
- ✤ Goal:

- C
- \succ try to find out all the attack structures related to <u>*t*</u>
- \blacktriangleright find out as many <u>t</u>'s friends and distant neighbors as possible
- Attack Steps
 - 1) Mutual friends between <u>a</u> & <u>t</u>; <u>a</u>'s friends;
 - 2) Need a user set \underline{U}
 - a randomized user set (Type 1)
 - a community /group including <u>*t*</u>(Type 2)
 - 3) Find out the attack structures and query mutual friends based on them





Exploration Attack

- Attacker (\underline{a}) has no specific target
- ✤ Goal
 - explore users who can be compromised
 - ➢ find out only one friend or distant neighbor of a user
- ✤ Attack Steps:
 - 1) Need a target set \underline{T}
 - a randomized user set (Type 1)
 - a community / group including <u>a (Type 2)</u>
 - 2) For each \underline{t} in \underline{T} ,
 - if there is one attack structure involving <u>a</u> & <u>t</u>
 - <u>t</u> can be compromised





- 1. Launch the exploration attack using attacker's community
- 2. Assume the attacker is interested in attacking t_1 and t_2 , he chooses them as specific targets
- 3. Launch specific-target attacks for t_1 and t_2 using t_1 's community and a user set about t_2



Specific-target Attacks

	Randomize d Group with 100 Users	Target's Group
Average Exposed Friends of a Target	2.6 (10.2%)	12.4 (48.9%)
Average Exposed Distant Neighbors of a Target	24.7	42.0



Exploration Attacks

	Randomized Group with 100 Users	Attacker's Group	
Average Compromised Users	12.3 (12.3%)	52.65 (61.0%)	



Hybrid Attacks

- Exploration + Specific-target
 - ✤ Launch Exploration attacks using the attacker's groups
 - Choose two specific targets who have the most mutual friends with the attacker from the Exploration attack results
 - Launch Specific-target attacks for selected targets using targets' groups

	Results of Hybrid Attacks	Results of Specific target Attack usin	
Average Exposed Friends of a Target	19.4 (73.2%)	~ 12 (49%)	
Average Exposed Distant Neighbors of a Target	48.3	~ 42	



Defense Approaches

• Reason

*no restriction for querying mutual friends

- Defense approaches
 - ✦Hide user profile

Access control to query mutual friends





Friendship Identification & Inference Attack

- Users' Privacy Settings for Friend Lists
 - Private
 - Friends w/o an excluding list
 - Public



Edit Privacy

Friend List



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Inconsistent Policies





Inconsistent Preferences Example -1





Inconsistent Preferences Example -2





Attack Definition

A social graph G(V, E); An adversary $b \in V$; A target $t \in V$

Assumption on Target t's Privacy Setting:

t defines a policy that does not authorize *b* to see *F(t)*

Assumptions on Adversary's Initial Knowledge:

- * **b**'s initial attack knowledge $K(b) = (V_{kb'}, E_{kb})$ is constructed based on friend lists visible to him
- * *t* is included in V_{kb}

Privacy Attack:

t is a victim of an *Fll* attack launched by *b* if *b* can identify and correctly infer at least *e* friends of *t*'s friends based on the proposed **random walk based link predictions** on *K(b)*.

No Additional Action



Attack Schemes

• One attacker node & one target



- Scheme 1: Adversary chooses a number of users, who are the most likely to be friends of a target, at one time
- Scheme 2: Adversary choose only one user, who is the most likely to be friend of a target at one time; add such a friendship link to the network and launch the attack again...
- Scheme 3: Adversary first attacks other users, who are close to the target; add identified and inferred friendship links to the network; then attack the target



Attack Schemes (cont.)

• Multiple attacker nodes & one target

Combine the attack knowledge (segments of the network) from different attacker nodes to be a more completed segment of the network

• Topology of the entire social network (multiple attacker nodes & multiple targets)

Attack the most vulnerable targets first





Results of attack scheme 1 in the three datasets

	Relationship Between an Attacker Node & a Target	D1	D2	D3
Average True Positive for Attacker Nodes (Average RC)	Friend	7.82 (78.2%)	4.71 (47.1%)	5.48 (54.8%)
	2-distant Neighbor	5.78 (57.8%)	2.85 (28.5%)	3.25 (32.5%)
	More than 2-distant Neighbor	4.03 (40.3%)	3.13 (31.3%)	3.19 (31.9%)
Average False Positive for Attacker Nodes (Average RIC)	Friend	2.18 (21.8%)	5.29 (52.9%)	4.52 (45.2%)
	2-distant Neighbor	4.22 (42.2%)	7.15 (71.5%)	6.75 (67.5%)
	More than 2-distant Neighbor	5.97 (59.7%)	6.87 (68.7%)	6.81 (68.1%)



Results of attack schemes in the three datasets

	Attack Scheme	D1	D2	D3
Average True Positive for Attacker Nodes	Attack Scheme 1	5.88	3.56	3.97
	Attack Scheme 2	5.92	3.78	4.22
	Attack Scheme 3	6.91	4.85	5.32



Attack results using multiple attacker nodes in D3

- M1: one friend, one 2-distant neighbor, one more than 2-distant neighbor
- M2: three friend, one 2-distant neighbor, one more than 2-distant neighbor
- M3: three friend, five 2-distant neighbor, five more than 2-distant neighbor
- ↔ M4: five friend, five 2-distant neighbor, five more than 2-distant neighbor

- Average True Positive for Targets

---- Average False Positive for Targets





Results from *FII* **attacks on the topology of the entire network in D3**

Each attack is repeated in 10 times

	After Launching 1K FII Attacks	After Launching 5K FII Attacks	After Launching 10K FII Attacks	After Launching 20K FII Attacks
Average Correctly Inferred Friendship Links (Percentage of All Friendship Links)	4113.4 (7.76%)	17948.3 (33.88%)	32231.4 (60.84%)	37891.2 (71.53%)
Average Incorrectly Friendship Links	4528.6	13314.4	21653.2	25472.3



Defense Approaches



- Squicciarini et al. -> voting algorithm & game theory
- Hu et al. -> Label Privacy Level, minimize privacy risk & sharing loss



Issues Related to Users' Posts

- Photos
 - A photo includes multiple individuals
 - One of them posts it in his/her wall
 - Other may be upset
- Check-ins (LBSNs) [2]
 - A user exposes where and when he is
 - A user exposes where his lives
 - A user's friend or other people expose the user's location related information
- Existing Access Control mechanisms cannot address all of these problems [5]



Other Issues – Email Address as Identity [7]

- Too many online systems adopt a user's email address as the user's identity
- Caused and causing many threats
 - Email address is not considered to be a private information
 - Easy to guess a user's identity in a online system
 - More vulnerable for online password cracking
 - Share the same passwords
 - Avoid the limits of fail login times
 - Cracking one email address = Cracking related online accounts associated with this email address



Social Media Landscape 2013



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Questions?

Thank You!

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