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# IS 2610: Data Structures

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Discuss HW 2 problems  
Binary Tree (continued)  
Introduction to Sorting

Feb 9, 2004

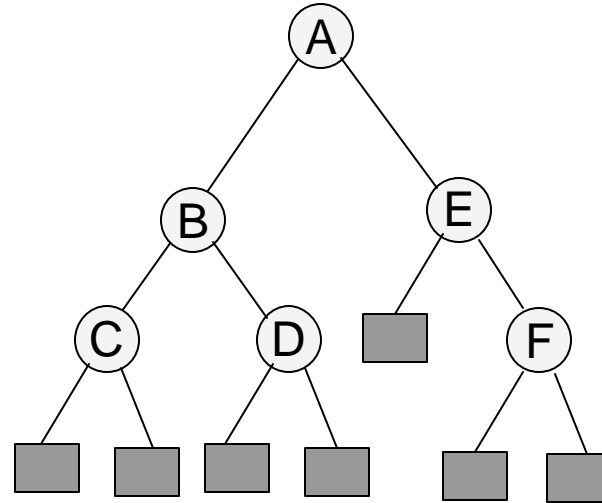
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# Binary tree

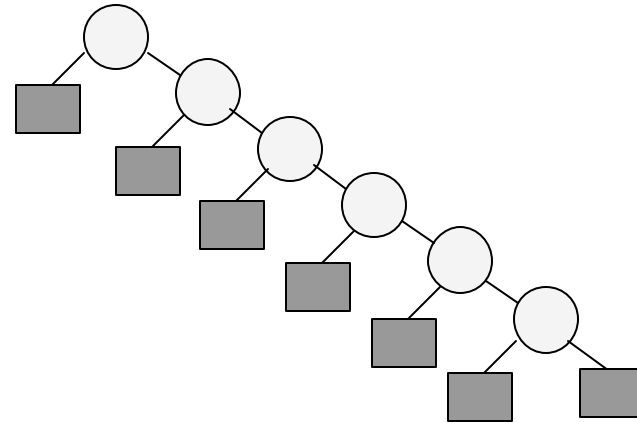
- A binary tree with  $N$  internal nodes has  $2N$  links
    - $N-1$  to internal nodes
      - Each internal node except root has a unique parent
      - Every edge connects to its parent
    - $N+1$  to external nodes
  - Level, height, path
    - Level of a node is  $1 +$  level of parent (Root is at level 0)
    - Height is the maximum of the levels of the tree's nodes
    - Path length is the sum of the levels of all the tree's nodes
    - Internal path length is the sum of the levels of all the internal nodes
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# Examples

- Level of D ?
- Height of tree?
- Internal length?
- External length?



- Height of tree?
- Internal length?
- External length?



# Binary Tree

- External path length of any binary tree with  $N$  internodes is  $2N$  greater than the internal path length
  - The height of a binary tree with  $N$  internal nodes is at least  $\lg N$  and at most  $N-1$ 
    - Worst case is a degenerate tree:  $N-1$
    - Best case: balanced tree with  $2^i$  nodes at level  $i$ .
      - Hence for height:  $2^{h-1} < N+1 = 2^h$  – hence  $h$  is the height
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# Binary Tree

- Internal path length of a binary tree with  $N$  internal nodes is at least  $N \lg (N/4)$  and at most  $N(N-1)/2$ 
    - Worst case :  $N(N-1)/2$
    - Best case:  $(N+1)$  external nodes at height no more than  $\lfloor \lg N \rfloor$ 
      - $(N+1) \lfloor \lg N \rfloor - 2N < N \lg (N/4)$
-

# Tree traversal (binary tree)

## ■ Preorder

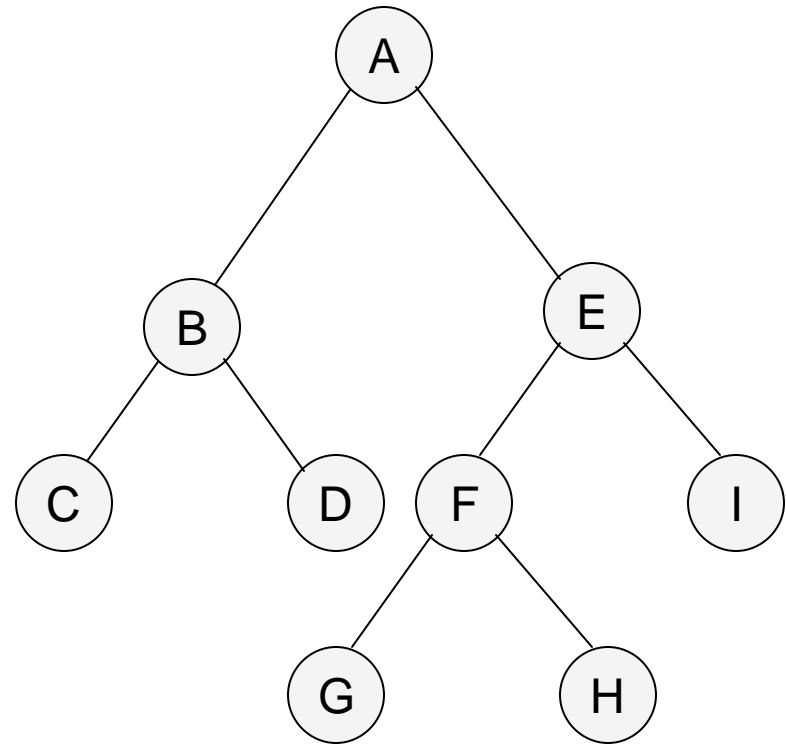
- Visit a node,
- Visit left subtree,
- Visit right subtree

## ■ Inorder

- Visit left subtree,
- Visit a node,
- Visit right subtree

## ■ Postorder

- Visit left subtree,
- Visit right subtree
- Visit a node



# Recursive/Nonrecursive Preorder

```
void traverse(link h, void (*visit)(link))
{
    If (h == NULL) return;
    (*visit)(h);
    traverse(h->l, visit);
    traverse(h->r, visit);
}
```

```
void traverse(link h, void (*visit)(link))
{
    STACKinit(max);
    STACKpush(h);
    while (!STACKempty())
    {
        (*visit)(h = STACKpop());
        if (h->r != NULL) STACKpush(h->r);
        if (h->l != NULL) STACKpush(h->l);
    }
}
```

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# Recursive binary tree algorithms

- Exercise on recursive algorithms:
  - Counting nodes
  - Finding height



# Sorting Algorithms: Selection sort

## ■ Basic idea

- Find smallest element and put in the first place
- Find next smallest and put in second place
- ..

## ■ Try for

2 6 3 1 5

```
#define key(A) (A)
#define less(A, B) (key(A) < key(B))
#define exch(A, B) { Item t = A; A = B; B = t; }

void selection (Item a[], int l, int r)
{
    int i, j;
    for (i = l; i < r; i++)
        { int min = i;
          for (j = i+1; j <= r; j++) {
              if (less(a[j], a[min]) min = j;
          }
          exch(a[i], a[min]);
        }
}
```

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# Sorting Algorithms: Selection sort

## ■ Recursive?

- Eliminate the outer loop
- Find the minimum and place it in the first place
- Make recursive call to sort the remaining parts of the array

## ■ Complexity

- $i$  goes from 1 to  $N-1$
  - There is one exchange per iteration; total  $N-1$
  - There is  $N-i$  comparisons per iteration
    - $(N-1) + (N-2) + \dots + 2 + 1 = N(N-1)/2$
-

# Sorting Algorithms: Bubble sort

## ■ Bubble sort

- Move through the elements exchanging adjacent pairs if the first one is larger than the second
- Try out !
  - 2 6 3 1 5

```
#define key(A) (A)
#define less(A, B) (key(A) < key(B))
#define exch(A, B) { Item t = A; A = B; B = t; }

void selection (Item a[], int l, int r)
{
    int i, j;
    for (i = l; i < r; i++)
        for (j = r; j > i; j--)
            if (less(a[j-1], a[j]) exch(a[j-1], a[j]);
}
```

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# Sorting Algorithms: Bubble sort

- Recursive?
  - Complexity
    - $i^{\text{th}}$  pass through the loop –  $(N-i)$  compare-and-exchange
    - Hence  $N(N-1)/2$  compare-and-exchange is the worst case
    - What would be the minimum number of exchanges?
-

# Sorting Algorithms: Insertion sort

## ■ Insertion sort

- “People”  
method

2 6 3 1 5

## ■ Complexity

- About  $N^2/4$
- *About half of  
the left array*

```
#define less(A, B) (key(A) < key(B))
#define exch(A, B) { Item t = A; A = B; B = t; }
#define compexch(A, B) if (less(B, A)) exch(A, B)

void insertion(Item a[], int l, int r) {
    int i;
    for (i = r; i > l; i--) compexch(a[i-1], a[i]);
    for (i = l+2; i <= r; i++) {
        int j = i; Item v = a[i];
        while (less(v, a[j-1])) {
            a[j] = a[j-1]; j--;
        }
        a[j] = v;
    }
}
```

# Sorting Algorithms: Shell sort

- Extend of insertion sort
  - Taking every  $h^{\text{th}}$  element will
  - Issues
    - Increment sequence?

$h = 13$

ASORTINGEXAMPLE

ASORTINGEXAMPLE

ASORTINGEXAMPLE

AEORTINGEXAMPLS

```
void shellsort(Item a[], int l, int r)
{ int i, j, h;
  for (h = 1; h <= (r-l); h = 3*h+1) ;
  for ( ; h > 0;)
    h = h/3;
    for (i = l+h; i <= r; i++)
      { int j = i; Item v = a[i];
        while (j >= l+h && less(v, a[j-h]))
          { a[j] = a[j-h]; j -= h; }
        a[j] = v;
      }
}
```

# Sorting Algorithms: Shell sort

$h = 4$

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...

```
void shellsort(Item a[], int l, int r)
{ int i, j, h;
  for (h = 1; h <= (r-l); h = 3*h+1) ;
  for ( ; h > 0;)
    h = h/3;
    for (i = l+h; i <= r; i++)
      { int j = i; Item v = a[i];
        while (j >= l+h && less(v, a[j-h]))
          { a[j] = a[j-h]; j -= h; }
        a[j] = v;
      }
}
```

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# Sorting Algorithms: Quick sort

- A divide and conquer algorithm
    - Partition an array into two parts
    - Sort the parts independently
    - Crux of the method is the partitioning process
      - Arranges the array to make the following three conditions hold
        - The element  $a[i]$  is in the final place in the array for some  $l$
        - None of the elements in  $a[l] \dots a[i-1]$  is greater than  $a[i]$
        - None of the elements in  $a[i+1] \dots a[r]$  is less than  $a[i]$
-



# Sorting Algorithms: Quick sort

```
int partition(Item a[], int l, int r);  
void quicksort(Item a[], int l, int r)  
{ int i;  
  if (r <= l) return;  
  i = partition(a, l, r);  
  quicksort(a, l, i-1);  
  quicksort(a, i+1, r);  
}
```

```
int partition(Item a[], int l, int r)  
{ int i = l-1, j = r; Item v = a[r];  
  for (;;)   
  {  
    while (less(a[++i], v)) ;  
    while (less(v, a[--j])) if (j == l) break;  
    if (i >= j) break;  
    exch(a[i], a[j]);  
  }  
  exch(a[i], a[r]);  
  return i;  
}
```

# Sorting Algorithms: Quick sort

7 20 5 1 3 11 8 10 2 9

7 20 5 1 3 11 8 10 2 9

7 2 5 1 3 11 8 10 20 9

7 2 5 1 3 11 8 10 20 9

7 2 5 1 3 8 11 10 20 9

$i$  |  $j$

7 2 5 1 3 8 9 10 20 11

$i$  |

```
int partition(Item a[], int l, int r)
{ int i = l-1, j = r; Item v = a[r];
  for (;;)
  {
    while (less(a[++i], v)) ;
    while (less(v, a[--j])) if (j == l) break;
    if (i >= j) break;
    exch(a[i], a[j]);
  }
  exch(a[i], a[r]);
  return i;
}
```